

# 平面上の点集合の平衡分割問題

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## Abstract

平面上に  $mq$  個の白点集合  $S$  と  $nq$  個の赤点集合  $T$  があり、 $S \cup T$  のどの 3 点も同一直線上にはないものとする。すると、 $S \cup T$  を下記の 2 つの条件を満たすような  $q$  個の集合  $P_1, P_2, \dots, P_q$  に分割できるという予想を与え、これが  $n = 1$  のとき成り立つことを示したが、今回は  $n = 2$  のときにもこの予想が成り立つことを証明する。またその証明から、所望の分割を得る  $O(N^2 \log N)$  のアルゴリズムもえられる、ただし  $N = mq + nq$  である。

- (1) すべての  $1 \leq i < j \leq q$  に対して、 $\text{conv}(P_i) \cap \text{conv}(P_j) = \phi$  である。
- (2) すべての  $1 \leq i \leq q$  に対して、 $|P_i \cap S| = n$  and  $|P_i \cap T| = m$  がなりたつ。

## Balanced partitions of two sets of points in the plane

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## Abstract

We prove the following theorem: Let  $n = 1$ ,  $m \geq 2$  and  $q \geq 1$  be integers and let  $S$  and  $T$  be two disjoint sets of points in the plane such that no three points of  $S \cup T$  are on the same line,  $|S| = nq$  and  $|T| = mq$ . Then  $S \cup T$  can be partitioned into  $q$  disjoint subsets  $P_1, P_2, \dots, P_q$  satisfying the following two conditions: (i)  $\text{conv}(P_i) \cap \text{conv}(P_j) = \phi$  for all  $1 \leq i < j \leq q$ ; and (ii)  $|P_i \cap S| = n$  and  $|P_i \cap T| = m$  for all  $1 \leq i \leq q$ .

We can obtain an  $O(N \log N)$  time algorithm for finding the above desired partition by our proof, where  $N = mq + nq$ . We also proved that the above theorem holds for  $n = 1$ , and give a conjecture which says that the above theorem holds for all integer  $n \geq 3$ . We don't give a complete proof of this theorem because of lack of pages, and a complete proof can be seen in [3].

**Key Works:** balanced partition, point set, plane,

# 1 Introduction

For a set  $P$  of points in the plane, we denote by  $\text{conv}(P)$  the *convex hull* of  $P$ , which is the smallest convex set containing  $P$ . In [2], we proved the following theorem.

**Theorem A** *Let  $m$  be a positive integer, and let  $S_1, S_2$  and  $T$  be three disjoint sets of points in the plane such that no three points of  $S_1 \cup S_2 \cup T$  are collinear (i.e., no three points of it are on the same line) and  $|T| = (m-1)|S_1| + m|S_2|$ . Put  $q = |S_1 \cup S_2|$ . Then  $S_1 \cup S_2 \cup T$  can be partitioned into  $q$  subsets  $P_1, P_2, \dots, P_q$  which satisfy the following three conditions: (i)  $\text{conv}(P_i) \cap \text{conv}(P_j) = \phi$  for all  $1 \leq i < j \leq q$ ; (ii)  $|P_i \cap (S_1 \cup S_2)| = 1$  for all  $1 \leq i \leq q$ ; and (iii)  $|P_i \cap T| = m-1$  if  $|P_i \cap S_1| = 1$ , and  $|P_i \cap T| = m$  if  $|P_i \cap S_2| = 1$ .*

In view of theorem A with  $S_1 = \emptyset$ , we gave the following conjecture in [2].

**Conjecture B** *Let  $m \geq 2, n \geq 2$  and  $q$  be positive integers. Let  $S$  and  $T$  be two disjoint sets of points in the plane such that no three points of  $S \cup T$  are collinear,  $|S| = nq$  and  $|T| = mq$ . Then  $S \cup T$  can be partitioned into  $q$  subsets  $P_1, P_2, \dots, P_q$  satisfying the following two conditions: (i)  $\text{conv}(P_i) \cap \text{conv}(P_j) = \phi$  for all  $1 \leq i < j \leq q$ ; and (ii)  $|P_i \cap S| = n$  and  $|P_i \cap T| = m$  for all  $1 \leq i \leq q$ .*

The above conjecture is true when  $q = 2$  because the conjecture with  $q = 2$  is equivalent to well-known discrete Ham Sandwich Theorem on the plane ([1] p.212). In this paper we show that the conjecture is true in the case of  $n = 2$ .

**Theorem 1** *Let  $m \geq 2$  and  $q \geq 1$  be integers and let  $S$  and  $T$  be two disjoint sets of points in the plane such that no three points of  $S \cup T$  are collinear,  $|S| = 2q$  and  $|T| = mq$ . Then  $S \cup T$  can be partitioned into  $q$  disjoint subsets  $P_1, P_2, \dots, P_q$  satisfying the following two conditions:*

- (i)  $\text{conv}(P_i) \cap \text{conv}(P_j) = \phi$  for all  $1 \leq i < j \leq q$ ; and
- (ii)  $|P_i \cap S| = 2$  and  $|P_i \cap T| = m$  for all  $1 \leq i \leq q$ .

Let us note that from the proof of the above theorem, we can obtain a polynomial time algorithm for finding such a partition given in the theorem.

## 2 Proof of Theorem 1

In this paper, we deal with only *directed lines* in order to define the right side of a line and the left side of it. Thus a *line* means a directed line. A line  $l$  dissects the plane into three pieces:  $l$  and two open half-planes  $R(l)$  and  $L(l)$ , where  $R(l)$  and  $L(l)$  denote the *open half-planes* which are on the right side and on the left side of  $l$ , respectively (see Figure 1). Let  $r_1$  and  $r_2$  be two rays emanating from the same point  $p$ . Then  $r_1 \cup r_2$  dissects the plane into three pieces:  $r_1 \cup r_2$  and two open regions  $R(r_1) \cap L(r_2)$  and  $L(r_1) \cap R(r_2)$ , where  $R(r_1) \cap L(r_2)$  denotes the open region which is on the right side of  $r_1$  and on the

left side of  $r_2$ , and  $L(r_1) \cap R(r_2)$  denotes the other open region (see Figure 1). Namely,  $R(r_1) \cap L(r_2)$  denotes the open region that is swept by the ray being rotated clockwise around  $p$  from  $r_1$  to  $r_2$ . If the internal angle  $\angle r_1 p r_2 = \angle r_1 r_2$  of  $R(r_1) \cap L(r_2)$  is less than  $\pi$ , then we call  $R(r_1) \cap L(r_2)$  the *wedge* defined by  $r_1$  and  $r_2$ , and denote it by  $\text{wdg}(r_1 p r_2)$  or  $\text{wdg}(r_2 p r_1)$ . Let us note that  $p \notin R(r_1) \cap L(r_2)$  and  $p \notin \text{wdg}(r_1 p r_2)$  since they are open regions and do not contain their boundaries.

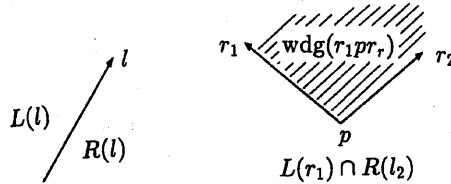


Figure 1: Open regions  $R(l)$ ,  $L(l)$  and  $L(r_1) \cap R(r_2)$  and a wedge  $\text{wdg}(r_1 p r_2) = R(r_1) \cap L(r_2)$ .

Let  $l_i$  be a line with suffix  $i$ , and  $p$  be a point on  $l_i$ . Then we denote by  $l_i^*$  the line which is obtained from  $l_i$  by reversing its direction. Moreover, we define the two rays  $r_i$  and  $r_i^*$  lying on the line  $l_i$  and having the same starting point  $p$  such that  $r_i$  has the same direction as  $l_i$  and  $r_i^*$  has the opposite direction of  $l_i$ . In particular,  $l_i = r_i \cup r_i^*$  (see Figure 2). Conversely, given a ray  $r_i$ , we can similarly define the ray  $r_i^*$ , whose direction is opposite to  $r_i$ , and the two lines  $l_i$  and  $l_i^*$ .

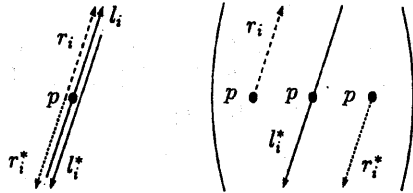


Figure 2: Lines  $l_i$  and  $l_i^*$ , and rays  $r_i$  and  $r_i^*$ .

For a region  $W$  in the plane, we define the integer-valued function  $f$  of  $W$  with respect to  $S$  and  $T$  by

$$f(W) := m|S \cap W| - 2|T \cap W|,$$

where  $S$  and  $T$  are the two disjoint sets of points in the plane given in Theorem 1. Hereafter  $f$  always denotes this function. A region  $W$  is said to be *balanced* if  $f(W) = 0$ . For example,  $\text{conv}(S \cup T)$  and  $\text{conv}(P_i)$  are balanced, where  $P_i$  is a subset of  $S \cup T$  given in Theorem 1. In order to prove Theorem 1, we need the following lemma.

**Lemma 2** *Let  $S$  and  $T$  be two disjoint sets of points in the plane given in Theorem 1. If there exist two lines  $l_1$  and  $l_2$  such that  $|R(l_1) \cap S| = |R(l_2) \cap S|$  and  $|R(l_1) \cap T| < |R(l_2) \cap T|$ , then for every integer  $i$ ,  $|R(l_1) \cap T| \leq i \leq |R(l_2) \cap T|$ , there exists a line  $l_3$  such that  $|R(l_3) \cap S| = |R(l_1) \cap S|$  and  $|R(l_3) \cap T| = i$ .*

**Proof** We first assume that  $R(l_1) \cap S = R(l_2) \cap S$  (see Figure 3). Then we can continuously move a line  $l$  from  $l_1$  to  $l_2$  in such a way that each line  $l$  passes through at most one point of  $T$  but no point of  $S$ . Then  $R(l) \cap S = R(l_1) \cap S$ , and the number  $|R(l) \cap T|$  changes  $\pm 1$  when  $l$  hits or passes a point of  $T$ . Therefore we can find the desired line  $l_3$ .

We next assume  $R(l_1) \cap S \neq R(l_2) \cap S$ . Consider two convex hulls  $\text{conv}(S \cap R(l_1))$  and  $\text{conv}(S \setminus R(l_1))$ . Then we can find two vertices  $x \in \text{conv}(S \cap R(l_1))$  and  $y \in \text{conv}(S \setminus R(l_1))$  such that a line  $l_4$  passing through  $x$  and  $y$  satisfies  $\angle l_2 l_4 < \angle l_2 l_1$  (see Figure 3). Let  $l'_4$  denote a line very close to  $l_4$  such that  $R(l'_4)$  contains  $x$  but not  $y$ , and  $l''_4$  denote a line very close to  $l_4$  such that  $R(l''_4)$  contains  $y$  but not  $x$ . We may assume that no point of  $(S \cup T) \setminus \{x, y\}$  lies between  $l'_4$  and  $l''_4$ . We can continuously move a line  $l$  from  $l_1$  to  $l'_4$  in such a way that  $l$  passes no point of  $S$  and the number  $|R(l) \cap T|$  changes  $\pm 1$ . Moreover it follows that  $R(l'_4) \cap T = R(l''_4) \cap T$  and  $R(l'_4) \cap S = ((R(l'_4) \cap S) \setminus \{x\}) \cup \{y\}$ .

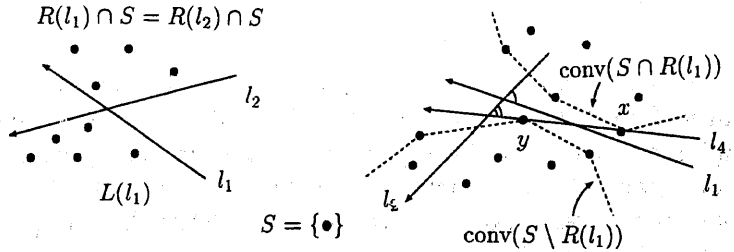


Figure 3: Two cases of the proof of Lemma 3.

Since  $|R(l_1) \cap S| = |R(l_2) \cap S|$ , by repeating this procedure we can obtain a line  $l_5$  possessing the property that  $R(l_5) \cap S = R(l_2) \cap S$  and that for every integer  $j$  between  $|R(l_1) \cap T|$  and  $|R(l_5) \cap T|$ , there exists a line  $l_6$  such that  $|R(l_6) \cap S| = |R(l_1) \cap S|$  and  $|R(l_6) \cap T| = j$ . Since  $l_2$  and  $l_5$  satisfy the condition of the previous case, the lemma is proved.  $\square$

**Proof of Theorem 1** We now prove Theorem 1 by induction on  $|S|$ . Unless otherwise stated, except when it moves, we always consider lines that pass through no points of  $S \cup T$ . We begin with the following Claim.

**Claim 3** For every integer  $i$ ,  $0 \leq i \leq q-1$ , there exist a line  $l$  that passes through two distinct points of  $S$  and satisfies  $|R(l) \cap S| = i$ .

**Proof** Let  $i$  be an integer such that  $0 \leq i \leq q-1$ . Let  $x$  be a vertex of  $\text{conv}(S)$  which lies on the bottom of  $\text{conv}(S)$ , and let  $l_1$  be the line that passes through  $x$  and the rightmost edge incident with  $x$  and goes upward. Then  $|R(l_1) \cap S| = 0$  and  $|L(l_1) \cap S| = 2q-2$ . By a suitable counterclockwise rotation of  $l_1$  around  $x$ , we can find a line  $l$  which passes through  $x$  and one more point of  $S$  and satisfies  $|R(l) \cap S| = i$  since no three points of  $S$  are collinear.  $\square$

**Claim 4** If  $q$  is even, then the theorem holds. Thus we may assume that  $q$  is odd. In particular, we can put  $q = 2k+1$ ,  $k \geq 1$ .

**Proof** Suppose that  $q$  is even. Then by Ham Sandwich Theorem, there exists a line  $l$  such that  $|R(l) \cap S| = |L(l) \cap S| = |S|/2$  and  $|R(l) \cap T| = |L(l) \cap T| = |T|/2$ , which imply that both  $R(l)$  and  $L(l)$  are balanced regions. By the inductive hypotheses on  $R(l)$  and on  $L(l)$ , we can obtain the required partition of  $S \cup T$ .  $\square$

By the same argument in the above proof, if there exists a line  $l$  such that  $f(R(l)) = 0$  and  $0 < |R(l) \cap S| < |S|$ , then both  $R(l)$  and  $L(l)$  are balanced, and thus we can obtain the desired partition of  $S \cup T$  by the inductive hypotheses on  $R(l)$  and on  $L(l)$ . Therefore we may assume that

$$f(R(l)) \neq 0 \quad \text{for every line } l \text{ with } 0 < |R(l) \cap S| < |S|. \quad (1)$$

We put

$$q = 2k + 1, \quad |S| = 4k + 2 \quad \text{and} \quad |T| = m(2k + 1).$$

**Claim 5** *We may assume that for every line  $l$  for which  $2 \leq |R(l) \cap S| = 2j \leq 2k$ , we have  $|R(l) \cap T| > mj$ , in particular,  $f(R(l)) < 0$  because otherwise the theorem holds.*

**Proof** If there exist two lines  $l_1$  and  $l_2$  such that  $2 \leq |R(l_1) \cap S| = |R(l_2) \cap S| = 2j \leq 2k$  and  $|R(l_1) \cap T| < mj < |R(l_2) \cap T|$ , then by Lemma 2 we can find a line  $l_3$  for which  $|R(l_3) \cap S| = 2j$  and  $|R(l_3) \cap T| = mj$ , which contradicts (1). Therefore the existence of a line  $l$  such that  $|R(l) \cap S| = 2j$  and  $|R(l) \cap T| > mj$  (or  $< mj$ ) is equivalent to the assertion that for every line  $l$  with  $|R(l) \cap S| = 2j$ , we have  $|R(l) \cap T| > mj$  (or  $< mj$ ). Thus it is enough to prove that for every  $1 \leq j \leq k$ , there exists a line  $l$  for which  $|R(l) \cap S| = 2j$  and  $|R(l) \cap T| > mj$ .

We prove this by induction on  $j$  from  $j = k$  to  $j = 1$ . By Claim 3, there exists a line  $l_4$  such that  $l_4$  passes through two points of  $S$  and  $|R(l_4) \cap S| = 2k$ , which implies  $|L(l_4) \cap S| = 2k$ . Since  $l_4$  does not pass through any point of  $T$  and by the equality  $|R(l_4) \cap S| = |L(l_4) \cap S|$ , we may assume that  $|R(l_4) \cap T| \geq |T|/2 > mk$ , and thus the statement holds when  $j = k$ .

Suppose that the claim holds for  $j + 1$  but does not for  $j$ , that is, assume that there exists a line  $l_1$  such that  $|R(l_1) \cap S| = 2j$  and  $|R(l_1) \cap T| < mj$ . Then for every line  $l_2$  with  $|R(l_2) \cap S| = 2j$ , we have  $|R(l_2) \cap T| < mj$ . By Claim 3, there exists a line  $l_3$  such that  $l_3$  passes through two points, say  $x$  and  $y$ , of  $S$  and  $|R(l_3) \cap S| = 2j$ . Since no three points of  $S \cup T$  are collinear, we can move  $l_3$  leftward very slightly so that the resulting line  $l_4$  satisfies that  $R(l_4) \cap S = (R(l_3) \cap S) \cup \{x, y\}$  and  $R(l_4) \cap T = R(l_3) \cap T$ . Thus  $|R(l_4) \cap S| = 2(j + 1)$  and  $|R(l_4) \cap T| < mj < m(j + 1)$ , which contradicts the fact that the claim holds for  $j + 1$ . Consequently the claim is proved.  $\square$

Let  $l_1$  be a line which passes through two points of  $S$ , say  $x$  and  $y$ , and satisfies  $R(l_1) \cap S = \emptyset$ . By a suitable rotation of the plane, we may assume that  $l_1$  lies horizontal and goes from right to left (see Figure 4). By considering a line  $l'_1$  lying very little below  $l_1$ , we have  $|R(l_1) \cap T| = |R(l'_1) \cap T| > m$  by  $|R(l'_1) \cap S| = 2$  and Claim 5. As it is easily seen, there exists a point  $z$  in  $L(l_1) \cap S$  such that letting  $l_2$  be the line passing through  $x$  and  $z$ ,  $|R(l_2) \cap S| = |L(l_2) \cap S| = 2k$ . We assume that  $l_2$  is directed from  $x$  to  $z$ , which

means that  $l_2$  goes downward (see Figure 4). Then by Claim 5, we have  $|R(l_2) \cap T| > mk$  and  $|L(l_2) \cap T| > mk$ . Let

$$a := |R(l_2) \cap T| - mk \quad \text{and} \quad b := |L(l_2) \cap T| - mk.$$

Since  $|R(l_2) \cap T| + |L(l_2) \cap T| = |T| = m(2k + 1)$ , we have

$$a > 0, \quad b > 0, \quad a + b = m, \quad f(R(l_2)) = -2a \quad \text{and} \quad f(L(l_2)) = -2b. \quad (2)$$

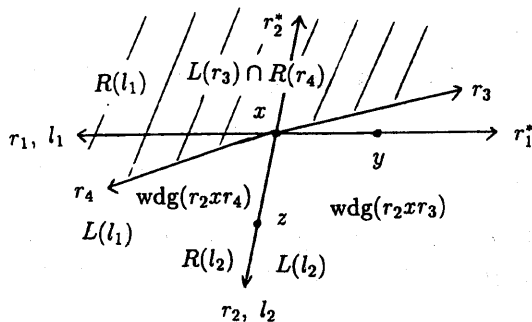


Figure 4: Lines  $l_1$ ,  $l_2$  and rays  $r_3$  and  $r_4$ .

Hereafter we consider rays emanating from  $x$ , and so, unless otherwise stated, a *ray* means such a ray.

**Claim 6** *Let  $r_2$  denote the ray lying on  $l_2$ . We may assume that there exists two rays  $r_3$  in  $L(l_2)$  and  $r_4$  in  $R(l_2)$  such that both wedges  $\text{wdg}(r_2xr_3)$  and  $\text{wdg}(r_2xr_4)$  are balanced and  $L(r_3) \cap R(r_4)$  contains exactly  $m$  points of  $T$  but no point of  $S$ . Of course,  $r_3$  must lie in  $\text{wdg}(r_1^*xr_2^*)$  and it may happen that  $r_4$  lies below  $r_1$  (see Figure 4).*

**Proof** We shall prove only the existence of  $r_3$  which satisfies that  $\text{wdg}(r_2xr_3)$  is balanced and  $\text{wdg}(r_3xr_2^*)$  contains exactly  $a$  points of  $T$  but no point of  $S$  because we can show the existence of  $r_4$  satisfying the similar conditions by the same argument, and the existence of these two rays implies Claim 6 by (2).

Recall that unless otherwise stated, we consider lines and rays which pass through no point of  $S \cup T$ . Note that an empty wedge  $\text{wdg}(r_2xr_2)$  has no point of  $S \cup T$  and is clearly balanced, that is,  $f(\text{wdg}(r_2xr_2)) = 0$ . We choose a ray  $r_3$  in  $L(l_2)$  so that

- (a)  $|\text{wdg}(r_2xr_3) \cap S|$  is even,
- (b)  $f(\text{wdg}(r_2xr_3)) \geq 0$ , and
- (c)  $|\text{wdg}(r_2xr_3) \cap (S \cup T)|$  is maximum subject to (a) and (b).

We begin with a observation that the value  $f(W)$  of a region  $W$  is always even when  $W$  contains even number of points in  $S$ , and that  $|L(l_2) \cap S| = 2k$ . We consider two cases.

**Case 1**  $\text{wdg}(r_3xr_2^*)$  contains at most  $m - 1$  points of  $S \cup T$ .

Since  $|\text{wdg}(r_3xr_2^*) \cap S|$  is even, if  $\text{wdg}(r_3xr_2^*)$  contains at least one point of  $S$ , then it contains at least two points of  $S$ , and so  $f(\text{wdg}(r_3xr_2^*)) \geq 2m - 2(m - 3) > 0$ . Hence  $f(L(l_2)) = f(\text{wdg}(r_2xr_3)) + f(\text{wdg}(r_3xr_2^*)) > 0$ , contradicting the fact that  $f(L(l_2)) <$

0. Therefore  $\text{wdg}(r_3xr_2^*)$  contains no point in  $S$ , and hence  $f(\text{wdg}(r_2xr_3)) = 0$  by the maximality (c) of  $|\text{wdg}(r_2xr_3) \cap (S \cup T)|$ . Consequently  $\text{wdg}(r_2xr_3)$  is balanced, and moreover  $\text{wdg}(r_3xr_2^*)$  contains exactly  $a$  points of  $T$  since  $f(L(l_2)) = f(\text{wdg}(r_2xr_3)) + f(\text{wdg}(r_3xr_2^*)) = 0 - 2|\text{wdg}(r_3xr_2^*) \cap T| = -2a$ .

**Case 2**  $\text{wdg}(r_3xr_2^*)$  contains at least  $m$  points of  $S \cup T$ .

In this case we shall prove that the theorem holds. Let  $r_5$  be a ray in  $\text{wdg}(r_3xr_2^*)$  such that  $\text{wdg}(r_3xr_5)$  contains exactly  $m$  points of  $S \cup T$  (see Figure 5). We distinguish two cases.

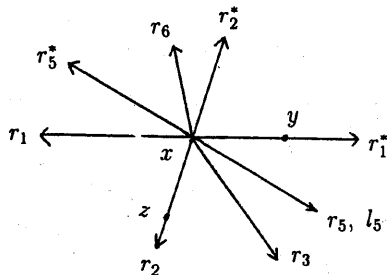


Figure 5: Lines  $l_5, l_6$  and rays  $r_5$  and  $r_6$ .

**Subcase 2.1**  $\text{wdg}(r_3xr_5)$  contains no point of  $S$ .

We omit the proof of this case.

**Subcase 2.2**  $\text{wdg}(r_3xr_5)$  contains at least one point of  $S$ .

We omit the proof of this case.  $\square$

We turn our attention to the proof of the theorem. Choose rays  $r_3$  and  $r_4$  according to Claim 6. Then it is obvious that  $(L(r_3) \cap R(r_4)) \cup \{x, z\}$  is balanced. In order to deal with a set of points of  $S \cup T$  contained in  $(L(r_3) \cap R(r_4)) \cup \{x, z\}$ , we consider a point  $x'$  on  $l_2$  and two rays  $r'_3$  and  $r'_4$  whose starting points are  $x'$ . First let  $x', r'_3$  and  $r'_4$  and are  $x, r_3$  and  $r_4$ , respectively. Then we continuously move  $x'$  on  $l_2$  toward the point  $z$  together with the rays  $r'_3$  and  $r'_4$  in such a way that both rays  $r'_3$  and  $r'_4$  pass through no point of  $S \cup T$ , and we stop moving if  $x'$  reaches  $z$  or at least one of rays  $r'_3$  and  $r'_4$  meets a point of  $S \cup T$  and  $x'$  cannot move more down. we consider two cases.

**Case 1**  $x'$  arrives at  $z$  (see Figure 6).

We omit the proof of this case.

**Case 2**  $x'$  cannot arrive at  $z$ .

In this case,  $x'$  stops above  $z$  since at least one of  $r'_3$  and  $r'_4$  meets a point of  $S \cup T$ . Without loss of generality, we may assume that  $r'_3$  meets a point of  $S \cup T$ . Then  $r'_3$  is tangent to  $\text{conv}(\text{wdg}(r_3xr_2^*) \cap T)$  at a vertex  $a$  and to  $\text{conv}(\text{wdg}(r_2xr_3) \cap (S \cup T))$  at a vertex  $b$  (see Figure 7).

Suppose that  $b$  is a point of  $S$ . Then  $(L(r'_3) \cap R(r'_4)) \cup \{a, b\}$ ,  $\text{wdg}(r_2xr'_3) \cup \{z\}$  and  $\text{wdg}(r_2xr'_4)$  are balanced regions, where  $L(r'_3) \cap R(r'_4) \ni x$  and  $\text{wdg}(r_2xr'_3) \cup \{z\} \not\ni b$ .

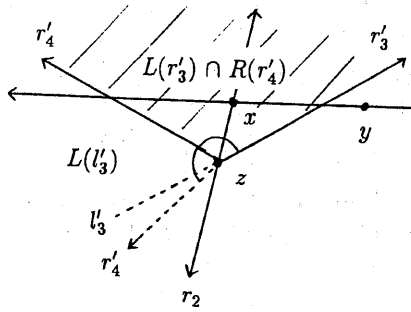


Figure 6:  $x'$  arrives at  $z$ .

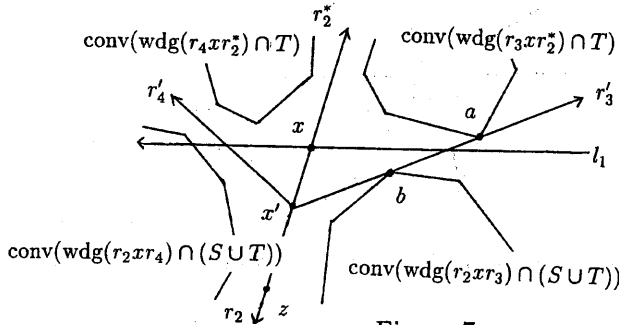


Figure 7:

Moreover, we can show that  $L(r_3') \cap R(r_4')$  is convex, (i.e.,  $\angle r_3' x' r_4' \leq \pi$ ) by applying Claim 5 to a line  $l_3'$  as in Case 1. Hence we can get the desired partition of  $S \cup T$  by inductive hypotheses. Therefore we may assume that  $b$  is a point of  $T$ .

It is clear that  $a$  is a point of  $T$ . Then we move  $x'$  very little down and define a new ray  $r_3'$  to be a ray that is very close to an old  $r_3'$  and passes through below  $b$  and above  $a$ . Then a new region  $L(r_3') \cap R(r_4')$  contains exactly  $m$  points of  $T$  and both new regions,  $\text{wdg}(r_2 x' r_3')$  and  $\text{wdg}(r_2 x' r_4')$  are balanced. Thus we can move  $x'$  toward to  $z$  again. By repeating this procedure, we can get the desired partition. Consequently the proof is complete.  $\square$

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