

Time Series Database Dedicated for a Computer Security Incident Response

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1. Introduction

In order to avoid data breach, it is important to quickly and accurately identify and confine a suspicious host when a computer security incident happens. When an external organization alerts a suspicious host, an IP address of the host is given. In order to identify the suspicious host, ones may then search for the IP address in related logs. Searching time of logs is then important to shorten a delay to identify a host. To this end, there have been already several logging systems such as fluentd[1][2], kibana[3], and splunk[4] that are based upon text messages of syslog. These existing logging systems are, however, not so efficient. For example, we have experienced that fluentd caused high CPU usages and lost log messages. In addition, these existing logging systems are not dedicated for a computer security incident. They, therefore, need more time to search for required logging messages for an incident.

This paper presents an idea of the novel fast logging database which is dedicated for a computer security incident response. The fast logging system tries to:

1. Scalable logging database with smaller storage
2. Fast search especially for recent logging messages

This rest of this paper is organized as follows. Section 2 presents an overview of the proposed fast logging database. Section 3 then presents basic key ideas of the fast logging database. Section 4 finally concludes this paper.

2. Overview of Fast Logging Database

Figure 1 depicts the overview of the

proposed fast logging database. The fast logging database consists of multiple database servers. All multiple database servers have the same IP address for IP anycasting. Network equipment or other servers (e.g., Web server, mail server and so on) send logging messages to the fast logging database using syslog protocol. The logging messages are then stored into one of database servers.

When one, say a CSIRT member, searches for logging messages, one sends a search request to one of database servers. The database server receiving a request forward the request to the other database servers. All database servers then send responses back to one who sends a search request.

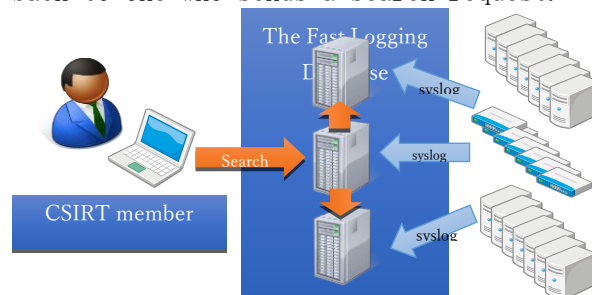


Figure 1 The Fast Logging Database

3. Fast Logging Database

The key ideas of the fast logging database can be summarized as follows:

1. Binary Based Key Value Store

Existing logging systems are basically based upon text messages while logging messages of network or security equipment usually are in pre-defined text format. Existing logging systems, therefore, have overhead to handle text messages. The fast logging database then stores binary values only in a record in a table, and text messages are indexed in another table.

2. Time Series Database (TSDB)

Since each logging message must have a timestamp, the fast logging database always stores the timestamp as a primary key. All records are basically stored in ascending order of timestamps. Some records are, however, not strictly in ascending order for lock-free operation described later.

3. Fixed Record Length

In order to improve search performance, the fast logging database has the same record length for a table as same as recent Relational Database Management System (RDBMS).

4. Timestamp Index

Regarding searching a record, a timestamp is usually specified for a computer security incident. In order to improve searching speed, a location of a record at a timestamp is indexed. Since the number of logging messages may depend upon daytime or night, timestamp index improves searching a record of a specified timestamp.

5. Logging Message Normalization

Logging message format in the fast logging database are automatically normalized.

Because a logging message is usually output using *printf* functions, the fast logging database indexes a message format. A Logging message is then stored as a tuple of a message format index and variable values, i.e., variable arguments of *printf*.

6. Lock-Free Clustering Support

The fast logging database does not strictly consider an order of a logging message. Timestamps in records are, therefore, not always in ascending order. This nature may delay finishing all search in order to make sure that the all log messages in specified timestamp in search are examined. This nature, however, makes insertion and lookup operations lock-free. Lock-free clustering can be then archived.

7. Recent in Memory and Old in Disk

When a computer security incident

happens and quick response is necessary, recent logging messages are searched in most cases. Older messages are not required to be fast to be searched because its search itself is enough delayed already. The fast logging database then always keeps recent logging messages in memory as much as possible, and write the messages to a disk if possible or all memory is consumed. Even after older messages are searched once, these older messages are not in memory in order to prioritize quick response for a recent computer security incident.

4. Concluding Remarks

This paper has presented the novel fast logging database that is dedicated for a computer security incident response. We have been implementing key ideas presented in this paper, and the evaluation in comparison with existing database systems[1][2][3][4][5] and other researches is future work.

5. References

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- [5] Y. Tubouchi, A. Wakisaka, K. Hamada, M. Matsuki, H. Abe and R. Matsumoto, “HeteroTSDB: A Time Series Database Architecture for Automatically Tiering on Heterogeneous Key-Value Stores,” Proc. Internet and Operation Technology Symposium (IOTS) 2018, pp.7-15, December 2018, Tottori, Japan.