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# Representation of SEBP Orientation by PLCP Orientation

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## 1. Introduction

Given a set of d+1 affinely independent points, how difficult it is to find the smallest d-dimensional ball enclosing all points? We call the problem the *Smallest enclosing ball of points* (SEB). This problem can be solved in linear time w.r.t. the number of points n with a fixed dimension of space [7, 10, 13]. In, these algorithms, however, the constant factor is exponential of dimension of space, hence a subexponential time algorithm algorithm was developed by Matoušek et al. [6]. Fischer and Gärtner gives  $O(d^3 \cdot (1.438)^d)$  algorithm [2] to SEBB using the unique sink orientations model, which was introduced by Szabó and Welzl [12].

This paper targets a special case of an SEB, which computes the smallest (d-1)-dimensional ball with a fixed one point on its boundary that encloses the other d points. We call this problem the smallest enclosing ball of points problem with a fixed point on the boundary (SEBP). An SEBP is also solved by Fischer and Gärtner's algorithm [2] as well as an SEB. We focus on the behavior of Fischer and Gärtner's algorithm to solve an SEBP. In the case of a ddimensional SEBP, the behavior is represented by a unique sink orientation on the graph  $G(C_d)$  of d-dimensional cube  $C_d$  [2, 12], called an SEBP-orientation. An orientation on  $G(C_d)$  is called a unique sink orientation (USO) of  $C_d$  if the induced orientation on every face F of  $C_d$  has a unique sink (and consequently has a unique source). A USO of a cube was introduced by Szabó and Welzl [12] as a framework to represent the behavior of some algorithms of mathematical programs. An outstanding example is obtained via a linear program (LP) solved by a simplex method. Let us consider an LP whose feasible region is combinatorially equivalent to  $C_d$ . If its linear function is in general position, every point of the feasible region has a different function value, hence every edge can be oriented from the endpoint with a smaller function value to a larger function value. The output orientation on  $G(C_d)$  is certainly a USO of  $C_d$ , called an LP-cube. The same is true of a linear complementarity problem (LCP), which is a unified framework for the two important optimization problems, linear programs and quadratic programs. If an LCP is solved by a series of pivoting algorithms, called Bard-type algorithms [1], we call the bahavior the LCP-orientation. Particularly if a given matrix of a LCP is a P-matrix whose all principal minors are positive, Stickney and Watson [11] proved the LCP-orientation is also a USO of  $C_d$ , called a PLCP-cube.

SEBP-orientations, which is a target of this paper, LP-cubes and PLCP-cubes have much importance as combinatorial properties of mathematical programs. In the case of LP-cubes and PLCP-cubes, other properties other than a USO were already discovered. It is easy to find that every LP-cube is acyclic, but as already observed by Stickney and Watson [11] there exists a PLCP-cube which has a directed cycle. On the other hand, LP-cubes and PLCP-cubes have the common property, called *Holt-Klee condition* [5]. Holt and Klee [5] proved that every LP-cube satisfies the Holt-Klee condition. Moreover, Gärtner, Morris and Rüst [4] recently resolved that PLCP-cubes also satisfies the Holt-Klee condition.

Differently from LP-cubes and PLCP-cubes, no properties of SEBP-orientations other than a USO are directly revealed, but Szabó et al. [12, 3] studied the formulation of SEB as finding sink in a cube graph, which is called an SEB-orientation. In the case of a d-SEB, the SEB-orientation is a USO of a  $C_{d+1}$  such that the induced orientation on the facet not including the global source is a SEBP-orientation. Gärtner, Miyazawa et al. [3] showed that an SEB-orientation has a directed cycle in sufficiently large dimension. However, SEBP-orientations including SEB-orientations, LP-cubes and PLCP-cubes have been not also characterized yet. All that is certain about three orientations is the inclusion relation between LP-cubes and PLCP-cubes, i.e., every LP-cube is a PLCP-cube, which was proved by Morris [8].

In this paper, we focus on another inclusion relation between SEBP-orientations and PLCP-cubes to show that every SEBP-orientation is a PLCP-cube.

#### 2. SEBP-orientation

**Definition 1** (A Smallest Enclosing Ball of Points with fixed point on boundary). Given P of d+1 affinely independent points in d-space. Choose one point v as origin, and the position of the other points are given as vectors  $\mathbf{x}_1, \mathbf{x}_2, \ldots, \mathbf{x}_d$ . Smallest enclosing ball of points with fixed point on boundary (SEBP) is to compute the smallest ball with v on its boundary which encloses all points in P whose dimension is d.

**Definition 2** (A SEBP-orientation). Consider the d dimensional cube whose vertex set is labeled  $\{Q \mid Q \subseteq P\}$ ,

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and edge set is  $\{\{Q \setminus \{x\}, Q\} \mid Q \subseteq P, x \in Q\}$ .  $\beta(X)$  denotes smallest ball that have all elements of point set  $X \subseteq \{\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_d\}$  and fixed point v on its boundary. Let the edges between  $\{Q \setminus \{x\}\}$  and  $\{Q\}$  be directed toward  $\{Q \setminus \{x\}\}$  iff  $x \in \beta(Q \setminus \{x\})$ .

The output orientation is a SEBP-orientation whose dimension is d.

If the sink of this orientation is  $S \subseteq P$ , the smallest enclosing ball of given points is  $\beta(S)$ , and it is uniquely determined.

**Definition 3** (A Linear complementarity problem). Given  $M \in \mathbb{R}^{n \times n}$  and  $q \in \mathbb{R}^n$ , find vectors  $w = (w_1, w_2, \dots, w_n)$ , and  $z = (z_1, z_2, \dots, z_n)$  such that

$$w - Mz = q, w \ge 0, z \ge 0, w^{\mathsf{T}}z = 0.$$
 (1)

If all principle minor of M is positive, it is called LCP on P-matrices (PLCP).

**Definition 4** (A PLCP-orientation). Consider the d-dimensional cube whose vertex set is labeled  $\{X \mid X \subseteq [d] := \{1, 2, \ldots, d\}\}$ . Let X be a set of index basis and  $M'_X$  be a submatrix of M whose i-th column is  $M_i$  if  $i \in X$ , otherwise unit vector  $\mathbf{e_i}$  Let the edges between  $X \setminus \{x\}$  and X be directed toward  $X \setminus \{x\}$  to X iff  $(M_X q)_x > 0$ . The output orientation is a PLCP-orientation in  $\mathcal{R}^d$ .

If the sink of PLCP-orientation is S, the basis of the feasible solution of given P-linear complementarity problem is  $\{z_k, w_l \mid k \in S, l \in [d] \setminus S\}$ .

**Theorem 1.** A SEBP-orientation is represented by A PLCP orientation.

Theorem 1. implies Corollary 1.

**Corollary 1.** A SEBP-orientation satisfies the Holt-Klee condition.

We prove Theorem 1. by inductive method of dimension. First, we represent SEBP by PLCP.

**Lemma 1.** d-dimensional SEB can be formulated to LCP with following P-matrix M and q,

$$M = \begin{pmatrix} 2\mathbf{x}_{1}^{\top}\mathbf{x}_{1} & 2\mathbf{x}_{1}^{\top}\mathbf{x}_{2} & \dots & 2\mathbf{x}_{1}^{\top}\mathbf{x}_{d} \\ 2\mathbf{x}_{2}^{\top}\mathbf{x}_{1} & 2\mathbf{x}_{2}^{\top}\mathbf{x}_{2} & \dots & 2\mathbf{x}_{2}^{\top}\mathbf{x}_{d} \\ \vdots & \vdots & \ddots & \vdots \\ 2\mathbf{x}_{d}^{\top}\mathbf{x}_{1} & 2\mathbf{x}_{d}^{\top}\mathbf{x}_{2} & \dots & 2\mathbf{x}_{d}^{\top}\mathbf{x}_{d} \end{pmatrix}, \qquad (2)$$

$$q = \begin{pmatrix} -\mathbf{x}_{1}^{\top}\mathbf{x}_{1} \\ -\mathbf{x}_{2}^{\top}\mathbf{x}_{2} \\ \vdots \\ -\mathbf{x}^{\top}\mathbf{x}_{d} \end{pmatrix} \qquad (3)$$

In the case d=2, it is easy to show every SEBP-orientation of 2-cube is PLCP-orientation.

Then we assume every SEBP-orientation is PLCP-orientation in  $d(\leq k-1)$ . In k-dimensional cube, proper face is PLCP-orientation because of the assumption.

**Lemma 2.** Let  $p_O$  be the outer point of k-1 points  $p_1, p_2, \ldots, p_{k-1}$ . In k-dimensional cube derived from  $p_1, p_2, \ldots, p_k$ , the value of k-th column of  $M_{\{1,\ldots,k-1\}}^{-1}$  is  $\mathbf{x}_k(2\mathbf{x}_O - \mathbf{x}_k)$ , where  $\mathbf{x}_O$  is position vector of the outer point of  $p_O$ .

**Lemma 3.** Given a point p whose position vector is  $\mathbf{x}_p$  and set of points P. Let  $p_O$  be the outer point of P and  $\mathbf{x}_O$  be its position vector. Then  $\mathbf{x}_p(2\mathbf{x}_O - \mathbf{x}_p) \geq 0$  iff  $p \in \beta(P)$ .

These two lemma imply that every SEBP-orientation is PLCP-orientation in the case d = k.

## 3. Conclusion

We show that every SEBP-orientation is PLCP-orientation. We also show SEB-orientation satisfies Holt-Klee condition [9] from Corollary 1. But the inclusion relation between SEB-orientations and other orientations has not been characterized yet, which is interesting open problem.

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