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S₆-FACTORIZATION ALGORITHM
OF COMPLETE BIPARTITE GRAPHS

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1. Introduction

Let S_6 be a *star* on 6 vertices and $K_{m,n}$ be a *complete bipartite graph* with partite sets V_1 and V_2 of m and n vertices each. A spanning subgraph F of $K_{m,n}$ is called an S_6 -factor if each component of F is isomorphic to S_6 . If $K_{m,n}$ is expressed as an edge-disjoint sum of S_6 -factors, then this sum is called an S_6 -factorization of $K_{m,n}$.

2. S_6 -factorization of $K_{m,n}$

Notation 1. r, t, b : number of S_6 -factors, number of S_6 -components of each S_6 -factor, and total number of S_6 -components, respectively, in an S_6 -factorization of $K_{m,n}$.

t_1 (t_2) : number of components whose centers are in V_1 (V_2), respectively, among t S_6 -components of each S_6 -factor.

$r_1(u)$ ($r_2(v)$) : number of components whose centers are all u (v) for any u (v) in V_1 (V_2), respectively, among b S_6 -components.

Lemma 1. If $K_{m,n}$ has an S_6 -factorization then (i) $b=mn/5$, (ii) $t=(m+n)/6$, (iii) $r=6mn/5(m+n)$, (iv) $t_1=(5n-m)/24$, (v) $t_2=(5m-n)/24$, (vi) $r_1=(5n-m)n/20(m+n)$, (vii) $r_2=(5m-n)m/20(m+n)$, (viii) $m \leq 5n$ and (ix) $n \leq 5m$.

Lemma 2. If $K_{m,n}$ has an S_6 -factorization, then $K_{sm,sn}$ has an S_6 -factorization for every positive integer s .

3. S_6 -factorization algorithm of $K_{m,n}$

Case (1) $m=5n$: From Lemma 2, $K_{5n,5n}$ has an S_6 -factorization since $K_{5,5}$ is just S_6 .

Case (2) $n=5m$: Obviously, $K_{m,5m}$ has an S_6 -factorization.

Case (3) $m < 5n$ and $n < 5m$: Let $x=(5n-m)/24$ and $y=(5m-n)/24$. And let $(x,4y)=d$, $x=dp$, $4y=dq$, where $(p,q)=1$.

Lemma 3. Let $(p,q)=1$ and s be an integer. Then

(I) when $q \equiv 1, 2, 3, 4 \pmod{5}$,

$$m=5(p+q)(5p+q)s, n=(25p+q)(5p+q)s \quad (5p+q \equiv 1, 3 \pmod{4})$$

$$\text{or } m=5(p+q)(5p+q)s/2, n=(25p+q)(5p+q)s/2 \quad (5p+q \equiv 2 \pmod{4})$$

$$\text{or } m=5(p+q)(5p+q)s/4, n=(25p+q)(5p+q)s/4 \quad (5p+q \equiv 0 \pmod{4})$$

(II) when $q=5q'$ and $q' \equiv 1, 2, 3, 4 \pmod{5}$,

$$m=5(p+5q')(p+q')s, n=5(5p+q')(p+q')s \quad (p+q' \equiv 1, 3 \pmod{4})$$

$$\text{or } m=5(p+5q')(p+q')s/2, n=5(5p+q')(p+q')s/2 \quad (p+q' \equiv 2 \pmod{4})$$

$$\text{or } m=5(p+5q')(p+q')s/4, n=5(5p+q')(p+q')s/4 \quad (p+q' \equiv 0 \pmod{4})$$

(III) when $q=25q''$,

$$m=(p+25q'')(p+5q'')s, n=5(p+q'')(p+5q'')s \quad (p+5q'' \equiv 1, 3 \pmod{4})$$

$$\text{or } m=(p+25q'')(p+5q'')s/2, n=5(p+q'')(p+5q'')s/2 \quad (p+5q'' \equiv 2 \pmod{4})$$

$$\text{or } m=(p+25q'')(p+5q'')s/4, n=5(p+q'')(p+5q'')s/4 \quad (p+5q'' \equiv 0 \pmod{4}).$$

Notation 2. Let A and B be two sequences of the same size such as

$$A: a_1, a_2, \dots, a_u$$

$$B: b_1, b_2, \dots, b_u.$$

If $b_i = a_i + c$ ($i=1, 2, \dots, u$), then we write $B=A+c$. If $b_i = ((a_i + c) \bmod w)$ ($i=1, 2, \dots, u$), then we write $B=A+c \bmod w$, where the residuals $a_i + c \bmod w$ are integers in the set $\{1, 2, \dots, w\}$.

Lemma 4. $(p, q) = 1$ and $q \equiv 1, 2, 3, 4 \pmod{5}$
 $m = 5(p+q)(5p+q)s$, $n = (25p+q)(5p+q)s$
 $\implies K_{m, n}$ has an S_6 -factorization.

Proof. (Algorithm I) Put $s=1$. Let $x=(5n-m)/24$, $y=(5m-n)/24$, $t=(m+n)/6$, $r=6mn/5(m+n)$. Then $x=5p(5p+q)$, $y=q(5p+q)$, $t=(5p+q)^2$, $r=(p+q)(25p+q)$. Let $r_m=p+q$, $r_n=25p+q$, $m_0=m/r_m=5(5p+q)$, $n_0=n/r_n=5p+q$.

Consider two sequences R and C of the same size $25(5p+q)$.

$R: 1, 1, 1, 1, 1, 2, 2, 2, 2, \dots, 5(5p+q), 5(5p+q), 5(5p+q), 5(5p+q), 5(5p+q)$

$C: 1, 2, \dots, 25(5p+q)-1, 25(5p+q)$.

Construct $R_i = R + 5(i-1)(5p+q)$ ($i=1, 2, \dots, p$).

Construct $C_i = (C + 5(i-1) \bmod 25(5p+q)) + 25(i-1)(5p+q)$ ($i=1, 2, \dots, p$).

Consider two sequences R' and C' of the same size $5(5p+q)$.

$R': r_1, r_2, \dots, r_{5(5p+q)}$, where $r_i = (i-1)p + 1 \bmod 5(5p+q)$ ($i=1, 2, \dots, 5(5p+q)$)

$C': c_1, c_2, \dots, c_{5(5p+q)}$, where $c_i = n - (i-1)q \bmod q(5p+q)$ ($i=1, 2, \dots, 5(5p+q)$).

Construct $R'_i = R' + 5(i-1)(5p+q) + 5p(5p+q)$ ($i=1, 2, \dots, q$).

Construct $C'_i = (C' - (i-1) \bmod q(5p+q)) + 25p(5p+q)$ ($i=1, 2, \dots, q$).

Consider two sequences I and J of the same size $5t$.

$I: R_1, R_2, \dots, R_p, R'_1, R'_2, \dots, R'_q$

$J: C_1, C_2, \dots, C_p, C'_1, C'_2, \dots, C'_q$.

Let i_k and j_k be the k -th element of I and J , respectively ($k=1, 2, \dots, 5t$). Join two vertices i_k in V_1 and j_k in V_2 with an edge (i_k, j_k) ($k=1, 2, \dots, 5t$). Construct a graph F with two vertex sets $\{i_k\}$ and $\{j_k\}$ and an edge set $\{(i_k, j_k)\}$. Then F is an S_6 -factor of $K_{m, n}$. This graph is called an S_6 -factor constructed with two sequences I and J .

Construct I_i such that $I_i = I + (i-1)m_0 \bmod m$ ($i=1, 2, \dots, r_m$).

Construct J_j such that $J_j = J + (j-1)n_0 \bmod n$ ($j=1, 2, \dots, r_n$).

Construct S_6 -factors $F_{i,j}$ with I_i and J_j ($i=1, 2, \dots, r_m; j=1, 2, \dots, r_n$). Then $F_{i,j}$ are edge-disjoint and their sum is an S_6 -factorization of $K_{m, n}$. By Lemma 2, $K_{m, n}$ has an S_6 -factorization for every positive integer s . \square

Lemma 5. $(p, q) = 1$ and $q = 5q'$ and $q' \equiv 1, 2, 3, 4 \pmod{5}$

$m = 5(p+5q')(p+q')s$, $n = 5(5p+q')(p+q')s$

$\implies K_{m, n}$ has an S_6 -factorization.

Lemma 6. $(p, q) = 1$ and $q = 25q''$

$m = (p+25q'')(p+5q'')s$, $n = 5(p+q'')(p+5q'')s$

$\implies K_{m, n}$ has an S_6 -factorization.

References

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