複数の直方体を折れる共通の展開図に関する研究

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本稿では,複数の異なる直方体を折れる共通の展開図を見つける問題を研究する. 2008年に二つの異なる直方体を折れる展開図が無限に存在することがすでに示されている.本稿ではまず三つの異なる直方体,具体的には大きさ $1 \times 1 \times 5$, $1 \times 2 \times 3$, $0 \times 1 \times 11$ の直方体が折れる展開図(直交多角形)を示す.直方体として体積が0のものを許してはいるものの,これは先行研究における未解決問題に対する解である. さらに,体積0の直方体を認めるならば,長い帯状の紙を使って,いくらでも多くの体積0の直方体が折れることを示す.次に,直交多角形以外の多角形で複数の異なる箱が折れる展開図が存在するかどうかを考える.従来の結果では,線が直交しているか45度を単位にするものしか考えられてこなかった.本稿では二つの直方体が折れる,より一般的な展開図が存在することを示す.

Common Developments of Several Different Orthogonal Boxes

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We investigate the problem of finding common developments that fold to plural incongruent orthogonal boxes. It was shown that there are infinitely many orthogonal polygons that fold to two incongruent orthogonal boxes in 2008. In this paper, we first show that there is an orthogonal polygon that fold to three boxes of size $1 \times 1 \times 5$, $1 \times 2 \times 3$, and $0 \times 1 \times 11$. Although we have to admit a box to have volume 0, this solves the open problem mentioned in literature. Moreover, once we admit that a box can be of volume 0, a long rectangular strip can be folded to an arbitrary number of boxes of volume 0. We next consider for finding common non-orthogonal developments that fold to plural incongruent orthogonal boxes. In literature, only orthogonal folding lines or with 45 degree lines were considered. In this paper, we show some polygons that can fold to two incongruent orthogonal boxes in more general directions.



図 1 キュビガミパズル. Fig. 1 Cubigami.

1. Introduction

Since Lubiw and O'Rourke posed the problem in 1996^{1} , polygons that can fold to a (convex) polyhedron have been investigated. In a book about geometric folding algorithms by Demaine and O'Rourke in 2007, many results about such polygons are given²⁾ [Chapter 25]. Such polygons have an application in the form of toys and puzzles. For example, the puzzle "cubigami" (Figure 1) is developed by Miller and Knuth, and it is a common development of all tetracubes except one (of surface area 16). One of the many interesting problems in this area is that whether there exists a polygon that folds to plural incongruent orthogonal boxes. Biedl et al. answered "yes" by finding two polygons that fold to two incongruent orthogonal boxes³⁾ (see $also^{2)}$ [Figure 25.53]). Later, Mitani and Uehara constructed infinite families of orthogonal polygons that fold to two incongruent orthogonal boxes.

First, we give an affirmative answer to this open problem, at least in some weak sense. That is, we give a polygon that can fold to three incongruent orthogonal boxes of size

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 $0 \times 1 \times 11$, $1 \times 1 \times 5$, and $1 \times 2 \times 3$. Note that one of the boxes is degenerate, as it has a side of length 0. Such a box is sometimes called a "doubly covered rectangle" (e.g.,⁵⁾). For boxes of positive volume, the existence of three boxes with a common unfolding is still open.

The polygon is found as a side effect of the enumeration of common developments of boxes of size $1 \times 1 \times 5$ and $1 \times 2 \times 3$. In the previous result by Mitani and Uehara⁴⁾, they randomly generated common developments of these boxes, and they estimated the number of common developments of these boxes at around 7000. However, they overestimated it since their algorithm did not exclude some symmetric cases. We enumerate all common developments of size $1 \times 1 \times 5$ and $1 \times 2 \times 3$, which can be found on a Web page^{*1}. As a result, the number of common developments of these boxes is 2263. Among 2263 developments, the development in Figure 2 is the only one that can fold to $0 \times 1 \times 11$.

Once we admit that a box can be a doubly covered rectangle, we have a new view of this problem since a doubly covered rectangle seems to be easier to construct than a box of positive volume. Indeed, we show that a sufficient long rectangular strip can be folded to an arbitrary number of doubly covered rectangles.

Next we turn to another approach to this topic. In an early draft by Biedl et al.³⁾, they showed a common development of two boxes of size $1 \times 2 \times 4$ and $\sqrt{2} \times \sqrt{2} \times 3\sqrt{2}$ (Figure 3). In the development, two folding ways to two boxes are not orthogonal. That is, the set of folding lines of a box intersect the other set of folding lines by 45 degrees. This development motivates us to the following problem: Is there any common development of two incongruent boxes such that two sets of folding lines intersect by an angle different from 45 or 90 degrees? We give an affirmative answer to this question.

2. Common orthogonal developments of boxes of size $1 \times 1 \times 5$ and $1 \times 2 \times 3$

For a positive integer S, we denote by P(S) the set of three integers a, b, c with $0 < a \le b \le c$ and ab + bc + ca = S, i.e., $P(S) = \{(a, b, c) \mid ab + bc + ca = S\}$. When



図 2 三つの異なる箱を折れる共通の展開図. (a) 大きさ 1 × 1 × 5 の箱を折るときの折り線. (b) 大きさ 1 × 2 × 3 の箱を折るときの折り線. (c) 大きさ 0 × 1 × 11 の箱を折るときの折り線.

Fig. 2 A common development of three different boxes. (a) Folding lines to make a 1 × 1 × 5 box.
(b) Folding lines to make a 1 × 2 × 3 box. (c) Folding lines to make a 0 × 1 × 11 box.

we only consider the case that folding lines are on the edges of unit squares, it is necessary to satisfy $|P(S)| \ge k$ to have a polygon of size 2S that can fold to k incongruent orthogonal boxes of positive volumes. The smallest S with $P(S) \ge 2$ is 11 and we have $P(11) = \{(1, 1, 5), (1, 2, 3)\}$. In this section, we concentrate at this special case. That is, we consider the developments that consist of 22 unit squares. Mitani and Uehara developed two randomized algorithms that try to find common developments of two different boxes⁴. Both algorithms essentially generate common developments randomly. Using the faster algorithm, they also estimated the number of common developments of the boxes of size $1 \times 1 \times 5$ and $1 \times 2 \times 3$ at around 7000. However, they overestimated it since their algorithm did not exclude some symmetric cases.

^{*1} http://www.jaist.ac.jp/~uehara/etc/origami/net/all-22.html

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図 3 Biedl たちによる二つの異なる箱の共通の展開図³⁾. (a) 大きさ $1 \times 2 \times 4$ の箱を折るときの折り線. (b) 大きさ $\sqrt{2} \times \sqrt{2} \times 3\sqrt{2}$ の箱を折るときの折り線.



We develop another algorithm that tries all common developments of these boxes. For a common development P of the boxes, let P' be a connected subset of P. That is, P' be a set of unit squares and it produces a connected simple polygon. Then, clearly, we can stick P' on these two boxes without overlap. We use the term *common partial development* of the boxes to denote such a smaller polygon. For example, one unit square is the common partial development of the boxes of surface area 1, and a rectangle of size 1×2 is the common development of them of surface area 2, and so on. Let L_i be the set of common partial developments of the boxes of surface area *i*. Then $|L_1| = |L_2| = 1$, and $|L_3| = 2$, and one of our main results is $|L_{22}| = 2263$. The outline of the first algorithm is described in Figure 4.

We implemented the algorithm and obtain all common developments in L_{22}^{*1} . One can find all of them at http://www.jaist.ac.jp/~uehara/etc/origami/net/all-22. html. All the values of L_i with $1 \le i \le 22$ are shown in Table 1. The first main theorem is as follows:

Theorem 1 The number of the common developments of boxes of size $1 \times 1 \times 5$ and $1 \times 2 \times 3$ into unions of unit squares is 2263.

let L_1 be a set of one unit square; for $i = 2, 3, 4, \dots, 22$ do

$$L_i \leftarrow \emptyset;$$

for each common partial development P in L_{i-1} do

for every polygon P^+ of size *i* obtained by attaching a unit square to P do check if P^+ is a common partial development, and add it into L_i if it is a new one;

end for

end for

end for

output L_{22} .

図 4 面積 22 の異なる箱が二つ折れるすべての多角形を出力するアルゴリズム.

Fig. 4 An algorithm that generates all common developments of two different boxes of area 22.

i	1	2	3	4	5	6	7	8	9	10	11	12
L_i	1	1	2	5	12	35	108	368	1283	4600	16388	57439
i-ominos	1	1	2	5	12	35	108	369	1285	4655	17073	63600

i	13	14	15	16	17	18
L_i	193383	604269	1632811	3469043	5182945	4917908
i-ominos	238591	901971	3426576	13079255	50107909	192622052

i	19	20	21	22	
L_i	2776413	882062	133037	2263	

表 1 大きさ $1 \times 1 \times 5$ の箱と大きさ $1 \times 2 \times 3$ の共通の部分展開図で,面積が $i(1 \le i \le 22)$ のものの個数(比較のため, $1 \le i \le 18$ に対しては i-オミノの個数も示した.)

Table 1 The number of common partial developments of two boxes $1 \times 1 \times 5$ and $1 \times 2 \times 3$ of surface area *i* with $1 \le i \le 22$. (For $1 \le i \le 18$, we give the number of *i*-ominos, for comparison.)

 $[\]star 1$ The first program with a naive implementation was too slow. We tuned it with many technical tricks, and now it outputs L_{22} in around 10 hours.

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3. Boxes including doubly-covered rectangles

3.1 Three boxes of surface area 22



Fig. 5 Tiling by the common development of three different boxes.

Among the 2263 developments in Theorem 1, there is only one development that gives an affirmative answer to the open problem in^{4} :

Theorem 2 There is a common development of three boxes of size $1 \times 1 \times 5$, $1 \times 2 \times 3$, and $0 \times 1 \times 11$. Moreover, the development is a polygon such that (1) it can fold to three boxes by orthogonal folding lines, and (2) it forms a tiling.

Proof. The development is depicted in Figure 2. It is easy to see that all folding lines in Figure 2(a)-(c) are orthogonal. The tiling is given in Figure 5.

In Theorem 2(1), one may complain that some folding lines are not on the edges of unit squares. Then, split each unit square into four unit squares. On the refined development for three boxes of surface area 88, we again have the claims in Theorem 2 for the boxes of size $2 \times 2 \times 10$, $2 \times 4 \times 6$, and $0 \times 2 \times 22$, and all folding lines are on the edges of unit squares.

3.2 A rectangular strip can be folded to an arbitrary number of doublycovered rectangles

Theorem 3 A rectangular $L \times 1$ paper (L > 1) can be folded into at least

$$2 + \lfloor L \rfloor$$

different doubly-covered rectangles in at least

$$1 + \left\lfloor \frac{L}{4} \right\rfloor + \left\lceil \frac{L}{4} \right\rceil + \left\lfloor L \right\rfloor$$

different ways.

Proof. Figure 6a shows how a long ribbon of width 1 can be wrapped by "twisting"



it around a rectangular strip. Here we show that we can obtain $\lfloor L \rfloor$ different doubly covered rectangles based on this way. First, we consider the points p_0, q_0, q_1, a, b, c , in Figure 6b). (Without loss of generality, we assume that $q_0b \ge q_1a$.) Let p_1 be the center of bc, and h_i is the point such that p_ih_i is a perpendicular of ab for i = 0, 1. We first observe that p_0a and bc are in parallel, the angles ap_0b and p_0bc are right angles, and p_0 is the center of q_0q_1 . Thus, careful analysis tells us that $\triangle q_0p_0b$, $\triangle h_0p_0b$, and $\triangle h_1p_1a$ are congruent. By symmetry, $\triangle q_1p_0a$, $\triangle h_0p_0a$, and $\triangle h_1p_1b$ are also congruent. Hence the points ap_0bp_1 form a rectangle. Therefore, the folding lines in Figure 6a) can be obtained by filling the rectangles like Figure 6b). Let k and w be the number of the rectangles and the length of the diagonal of the rectangle, respectively. Then, to obtain a feasible folding lines, we need $k \ge 1$, kw = L, and $w = ab \ge 1$. Therefore, for each $k = 1, 2, \ldots, \lfloor L \rfloor$, we can obtain a doubly covered rectangle of size p_0b and kp_0a .

In addition, we have the two ways of folding the ribbon in half along the long axis (leading to a $L \times \frac{1}{2}$ rectangle) or along the short axis (leading to a $(L/2) \times 1$ rectangle).

We next turn to another idea of folding. Figure 7a shows how a long ribbon of width 1 can be wrapped by "winding" it around a rectangular strip in such a way that the space between successive windings is equal to the width of the ribbon. By bending it backward at the end, as in Figure 7b–c, one obtains a doubly covered strip. Figure 7d

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shows the geometric construction: start with a right triangle ABC with the long side $d = BC = \cot \alpha + \tan \alpha$ on a long edge of the ribbon and the right angle A on the opposite edge. When the length L of the ribbon is an even multiple of d ($L = 2n \cdot d$), the folding will close into a doubly covered rectangle.



図 8 二重被覆長方形をリボンで折るための別の方法 . Fig. 8 A different way of folding a ribbon to a doubly-covered rectangle

The minimum possible value of d is 2. d changes continuously with α , and any value of d larger than 2 can be obtained. So n, the number of repetitions, can take all values between 1 and $n_{\max} := \lfloor L/4 \rfloor$. For each n in this range, one can form a right triangle ABC with hypotenuse d = L/(2n) and legs $\frac{1}{2}(\sqrt{d^2 + 2d} \pm \sqrt{d^2 - 2d})$. One can use the longer leg as the wrapping direction, as in Figure 7, or the shorter leg, as in Figure 8. This leads to doubly covered rectangles of dimensions $\left(n \cdot \frac{1}{2}(\sqrt{d^2 + 2d} + \sqrt{d^2 - 2d})\right) \times$ $\frac{1}{2}(\sqrt{d^2+2d}-\sqrt{d^2-2d}) \text{ and } \frac{1}{2}(\sqrt{d^2+2d}+\sqrt{d^2-2d}) \times \left(n \cdot \frac{1}{2}(\sqrt{d^2+2d}-\sqrt{d^2-2d})\right).$ For d = 2, the two possibilities coincide. So the total number of possibilities is $\lfloor L/4 \rfloor + \lceil L/4 \rceil - 1$. This equals $2\lfloor L/4 \rfloor$ except when L is a multiple of 4. In this case, we have to subtract 1 to compensate the overcounting for the case d = 2.

But we can see that each doubly covered rectangle by winding can be also obtained by twisting. Hence we obtain $2 + \lfloor L \rfloor$ different doubly covered rectangles in total.

4. Non-orthogonal polygons that fold to two incongruent boxes

Figure 9 shows a common unfolding of a $4 \times 4 \times 8$ box and a $\sqrt{10} \times 2\sqrt{10} \times 2\sqrt{10}$ box. It was obtained by solving an integer programming problem. The integer programming model formulates the problem of selecting a subset of 160 unit squares of the axis-aligned square grid underlying Figure 9, subject to the following constraints.

- (1) They should form a connected set in the plane.
- (2) When folded on the 4 × 4 × 8 box, every square of the surface is covered exactly once. (There are no overlaps.)
- (3) When folded on the $\sqrt{10} \times 2\sqrt{10} \times 2\sqrt{10}$ box, every part of the surface is covered exactly once. Note that the surface of the $\sqrt{10} \times 2\sqrt{10} \times 2\sqrt{10}$ box can be partitioned into 160 unit squares, which are however not aligned with the edges of the box. These squares result from folding the standard grid onto the box surface as shown in Figure 9. Some of these squares bend across an edge of the box.

The algorithm of Section 2 can be viewed as a systematic incremental way of finding all solutions to this problem.

The dimensions of the boxes were chosen as follows: A $1 \times 1 \times 2$ box has surface area 10, and a $1 \times 2 \times 2$ box has surface area 16. By scaling the first box with the factor 4 and the second box with the factor $\sqrt{10}$, we get two boxes with equal surface areas. A square lattice of side length $\sqrt{10}$ can be embedded on the standard integer grid by choosing the vector $\binom{1}{3}$ as a generating "unit vector".

The alignment of the two box unfoldings, with the symmetric layout of two "central" faces sharing two vertices, was fixed and was chosen by hand.

Figure 10 has been made from Figure 9 in an attempt to conceal the obvious folding directions. Further puzzles along these lines (for printing and cutting out) are given on

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a web page^{$\star 1$}.



図 9 二つの異なる箱が折れる共通の展開図.ある箱を折るための折り線は,他方の箱を折るための折り線とは直交しない.交わる角度は $\arctan 3 \approx 72^{\circ}$ である.

Fig. 9 A common development of two different boxes. The set of folding lines for one box intersect the other set by neither 90 nor 45 degrees, but at $\arctan 3 \approx 72^{\circ}$.

5. Concluding remarks

It is an open question if a polygon exists that can fold to three or more orthogonal boxes such that all of them have positive volume. We are exploring the possibility to find such examples by our integer programming model of Section 4. If we take the approach in Section 2, the smallest S with $|P(S)| \geq 3$ is given by $P(23) = \{(1,1,11), (1,2,7), (1,3,5)\}$. Thus we need to construct polygons of surface area 46, which is much bigger than 22.

In Section 3.2, we use three different ideas for folding a rectangular ribbon R to a



図 10 二つの異なる箱が折れる共通の展開図.図9の境界線を変更して作成.

Fig. 10 A common development of two different boxes. This has been obtained from Figure 9 by modifying the boundary.

doubly-covered rectangle. It would be interesting to classify *all* ways of folding ribbons into doubly-covered rectangles. In fact, we can generalize the ideas of "twisting" and "winding"; see Figure 11. These folding ways correspond to a kind of the billiard ball problem on a rectangular table. Hence, to specify all the folding ways in the figures, we have to find all pairs of relatively prime integers p and q with $pq = \lfloor cL \rfloor$ for c = 1, 1/4. The number of such pairs seems to be related to the maximal value of prime divisors of numbers in reduced residue system for $\lfloor cL \rfloor^{*2}$.

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^{*1} http://www.inf.fu-berlin.de/~rote/Software/folding-puzzles/

^{*2} http://oeis.org/A051265

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