## 複数の直方体を折れる共通の展開図に関する研究

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本稿では，複数の異なる直方体を折れる共通の展開図を見つける問題を研究する 2008 年に二つの異なる直方体を折れる展開図か無限に存在することがすでに示され ている。本稿ではまず三つの異なる直方体，具体的には大きさ $1 \times 1 \times 5,1 \times 2 \times 3$ $0 \times 1 \times 11$ の直方体が折れる展開図（直交多角形）を示す。直方体として体積が 0 の ものを許してはいるものの，これは先行研究における末解決問題に対する解である さらに，体積 0 の直方体を認めるならば，長い帯状の紙を使って，いくらでも多くの体積 0 の直方体が折れることを示す。次に，直交多角形以外の多角形で複数の異なる箱が折れる展開図が存在するかどうかを考える。従来の結果では，線が直交している か 45 度を単位にするものしか考えられてこなかった．本稿では二つの直方体か浙れ る，より一般的な展開図が存在することを示す。

## Common Developments of Several Different Orthogonal Boxes

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We investigate the problem of finding common developments that fold to plu－ ral incongruent orthogonal boxes．It was shown that there are infinitely many orthogonal polygons that fold to two incongruent orthogonal boxes in 2008．In this paper，we first show that there is an orthogonal polygon that fold to three boxes of size $1 \times 1 \times 5,1 \times 2 \times 3$ ，and $0 \times 1 \times 11$ ．Although we have to admit a box to have volume 0，this solves the open problem mentioned in literature． Moreover，once we admit that a box can be of volume 0，a long rectangular strip can be folded to an arbitrary number of boxes of volume 0 ．We next consider for finding common non－orthogonal developments that fold to plural incongruent orthogonal boxes．In literature，only orthogonal folding lines or with 45 degree lines were considered．In this paper，we show some polygons that can fold to two incongruent orthogonal boxes in more general directions．


図1 キュビガミパズル。 Fig． 1 Cubigami．

## 1．Introduction

Since Lubiw and O＇Rourke posed the problem in $1996^{1)}$ ，polygons that can fold to a （convex）polyhedron have been investigated．In a book about geometric folding algo－ rithms by Demaine and O＇Rourke in 2007，many results about such polygons are given ${ }^{2)}$ ［Chapter 25］．Such polygons have an application in the form of toys and puzzles．For example，the puzzle＂cubigami＂（Figure 1）is developed by Miller and Knuth，and it is a common development of all tetracubes except one（of surface area 16）．One of the many interesting problems in this area is that whether there exists a polygon that folds to plural incongruent orthogonal boxes．Biedl et al．answered＂yes＂by finding two polygons that fold to two incongruent orthogonal boxes ${ }^{3)}$（see also ${ }^{2)}$［Figure 25．53］）． Later，Mitani and Uehara constructed infinite families of orthogonal polygons that fold to two incongruent orthogonal boxes ${ }^{4)}$ ．However，it is open that whether there is a polygon that can fold to three or more boxes．

First，we give an affirmative answer to this open problem，at least in some weak sense． That is，we give a polygon that can fold to three incongruent orthogonal boxes of size

[^0]$0 \times 1 \times 11,1 \times 1 \times 5$ ，and $1 \times 2 \times 3$ ．Note that one of the boxes is degenerate，as it has a side of length 0 ．Such a box is sometimes called a＂doubly covered rectangle＂（e．g．，${ }^{5)}$ ）． For boxes of positive volume，the existence of three boxes with a common unfolding is still open．
The polygon is found as a side effect of the enumeration of common developments of boxes of size $1 \times 1 \times 5$ and $1 \times 2 \times 3$ ．In the previous result by Mitani and Uehara ${ }^{4)}$ ， they randomly generated common developments of these boxes，and they estimated the number of common developments of these boxes at around 7000．However，they overes－ timated it since their algorithm did not exclude some symmetric cases．We enumerate all common developments of boxes of size $1 \times 1 \times 5$ and $1 \times 2 \times 3$ ，which can be found on a Web page ${ }^{\star 1}$ ．As a result，the number of common developments of these boxes is 2263．Among 2263 developments，the development in Figure 2 is the only one that can fold to $0 \times 1 \times 11$ ．

Once we admit that a box can be a doubly covered rectangle，we have a new view of this problem since a doubly covered rectangle seems to be easier to construct than a box of positive volume．Indeed，we show that a sufficient long rectangular strip can be folded to an arbitrary number of doubly covered rectangles．

Next we turn to another approach to this topic．In an early draft by Biedl et al．${ }^{3)}$ ， they showed a common development of two boxes of size $1 \times 2 \times 4$ and $\sqrt{2} \times \sqrt{2} \times 3 \sqrt{2}$ （Figure 3）．In the development，two folding ways to two boxes are not orthogonal． That is，the set of folding lines of a box intersect the other set of folding lines by 45 degrees．This development motivates us to the following problem：Is there any common development of two incongruent boxes such that two sets of folding lines intersect by an angle different from 45 or 90 degrees？We give an affirmative answer to this question．

2．Common orthogonal developments of boxes of size $1 \times 1 \times 5$ and $1 \times 2 \times 3$

For a positive integer $S$ ，we denote by $P(S)$ the set of three integers $a, b, c$ with $0<a \leq b \leq c$ and $a b+b c+c a=S$ ，i．e．，$P(S)=\{(a, b, c) \mid a b+b c+c a=S\}$ ．When

[^1]

図 2 三つの異なる箱を折れる共通の展開図．（a）大きさ $1 \times 1 \times 5$ の箱を折るときの折り線．（b）大きさ $1 \times 2 \times 3$ の箱を折るときの折り線．（c）大きさ $0 \times 1 \times 11$ の箱を折るときの折り線
Fig． 2 A common development of three different boxes．（a）Folding lines to make a $1 \times 1 \times 5$ box． （b）Folding lines to make a $1 \times 2 \times 3$ box．（c）Folding lines to make a $0 \times 1 \times 11$ box．
we only consider the case that folding lines are on the edges of unit squares，it is nec－ essary to satisfy $|P(S)| \geq k$ to have a polygon of size $2 S$ that can fold to $k$ incongruent orthogonal boxes of positive volumes．The smallest $S$ with $P(S) \geq 2$ is 11 and we have $P(11)=\{(1,1,5),(1,2,3)\}$ ．In this section，we concentrate at this special case．That is，we consider the developments that consist of 22 unit squares．Mitani and Uehara developed two randomized algorithms that try to find common developments of two dif－ ferent boxes ${ }^{4)}$ ．Both algorithms essentially generate common developments randomly． Using the faster algorithm，they also estimated the number of common developments of the boxes of size $1 \times 1 \times 5$ and $1 \times 2 \times 3$ at around 7000 ．However，they overestimated it since their algorithm did not exclude some symmetric cases．


図 3 Biedl たちによる二つの異なる箱の共通の展開図 ${ }^{3)}$ ．（a）大きさ $1 \times 2 \times 4$ の箱を折るときの折り線．（b）大 きさ $\sqrt{2} \times \sqrt{2} \times 3 \sqrt{2}$ の箱を折るときの折り線 ．
Fig． 3 A common development of two different boxes by Biedl et al．${ }^{3}$ ．（a）Folding lines to make a $1 \times 2 \times 4$ box．（b）Folding lines to make a $\sqrt{2} \times \sqrt{2} \times 3 \sqrt{2}$ box．

We develop another algorithm that tries all common developments of these boxes． For a common development $P$ of the boxes，let $P^{\prime}$ be a connected subset of $P$ ．That is，$P^{\prime}$ be a set of unit squares and it produces a connected simple polygon．Then， clearly，we can stick $P^{\prime}$ on these two boxes without overlap．We use the term common partial development of the boxes to denote such a smaller polygon．For example，one unit square is the common partial development of the boxes of surface area 1 ，and a rectangle of size $1 \times 2$ is the common development of them of surface area 2 ，and so on． Let $L_{i}$ be the set of common partial developments of the boxes of surface area $i$ ．Then $\left|L_{1}\right|=\left|L_{2}\right|=1$ ，and $\left|L_{3}\right|=2$ ，and one of our main results is $\left|L_{22}\right|=2263$ ．The outline of the first algorithm is described in Figure 4.

We implemented the algorithm and obtain all common developments in $L_{22}{ }^{\star 1}$ ．One can find all of them at http：／／www．jaist．ac．jp／${ }^{\text {u uehara／etc／origami／net／all－22．}}$ html．All the values of $L_{i}$ with $1 \leq i \leq 22$ are shown in Table 1．The first main theorem is as follows：

Theorem 1 The number of the common developments of boxes of size $1 \times 1 \times 5$ and $1 \times 2 \times 3$ into unions of unit squares is 2263 ．

[^2]let $L_{1}$ be a set of one unit square；
for $i=2,3,4, \ldots, 22$ do
$$
L_{i} \leftarrow \emptyset
$$
for each common partial development $P$ in $L_{i-1}$ do
for every polygon $P^{+}$of size $i$ obtained by attaching a unit square to $P$ do check if $P^{+}$is a common partial development，and add it into $L_{i}$ if it is a new one；
end for
end for
end for
output $L_{22}$ ．
図4 面積 22 の異なる箱が二つ折れるすべての多角形を出力するアルゴリズム．
Fig． 4 An algorithm that generates all common developments of two different boxes of area 22 ．

| $i$ | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $L_{i}$ | 1 | 1 | 2 | 5 | 12 | 35 | 108 | 368 | 1283 | 4600 | 16388 | 57439 |
| $i$－ominos | 1 | 1 | 2 | 5 | 12 | 35 | 108 | 369 | 1285 | 4655 | 17073 | 63600 |


| $i$ | 13 | 14 | 15 | 16 | 17 | 18 |
| :---: | :---: | :---: | :---: | ---: | ---: | :---: |
| $L_{i}$ | 193383 | 604269 | 1632811 | 3469043 | 5182945 | 4917908 |
| $i$－ominos | 238591 | 901971 | 3426576 | 13079255 | 50107909 | 192622052 |


| $i$ | 19 | 20 | 21 | 22 |
| :---: | :---: | :---: | :---: | :---: |
| $L_{i}$ | 2776413 | 882062 | 133037 | 2263 |

表 1 大きさ $1 \times 1 \times 5$ の箱と大きさ $1 \times 2 \times 3$ の共通の部分展開図で，面積が $i(1 \leq i \leq 22)$ のものの個数．（比較のため， $1 \leq i \leq 18$ に対しては $i$－オミノの個数も示した．）
Table 1 The number of common partial developments of two boxes $1 \times 1 \times 5$ and $1 \times 2 \times 3$ of surface area $i$ with $1 \leq i \leq 22$ ．（For $1 \leq i \leq 18$ ，we give the number of $i$－ominos，for comparison．）

## 3．Boxes including doubly－covered rectangles

## 3．1 Three boxes of surface area 22



図 $\mathbf{5}$ 三つの異なる箱が折れる共通の展開図によるタイリング。
Fig． 5 Tiling by the common development of three different boxes．
different ways．
Proof．Figure 6a shows how a long ribbon of width 1 can be wrapped by＂twisting＂



図 6 二重被覆長方形をリボンで折るための別の方法．
Fig． 6 Another way of folding a ribbon to a doubly－covered rectangle．
it around a rectangular strip．Here we show that we can obtain $\lfloor L\rfloor$ different doubly covered rectangles based on this way．First，we consider the points $p_{0}, q_{0}, q_{1}, a, b, c$ ，in Figure 6b）．（Without loss of generality，we assume that $q_{0} b \geq q_{1} a$ ．）Let $p_{1}$ be the center of $b c$ ，and $h_{i}$ is the point such that $p_{i} h_{i}$ is a perpendicular of $a b$ for $i=0,1$ ．We first observe that $p_{0} a$ and $b c$ are in parallel，the angles $a p_{0} b$ and $p_{0} b c$ are right angles，and $p_{0}$ is the center of $q_{0} q_{1}$ ．Thus，careful analysis tells us that $\triangle q_{0} p_{0} b, \triangle h_{0} p_{0} b$ ，and $\triangle h_{1} p_{1} a$ are congruent．By symmetry，$\triangle q_{1} p_{0} a, \triangle h_{0} p_{0} a$ ，and $\triangle h_{1} p_{1} b$ are also congruent．Hence the points $a p_{0} b p_{1}$ form a rectangle．Therefore，the folding lines in Figure 6a）can be obtained by filling the rectangles like Figure 6 b ）．Let $k$ and $w$ be the number of the rectangles and the length of the diagonal of the rectangle，respectively．Then，to obtain a feasible folding lines，we need $k \geq 1, k w=L$ ，and $w=a b \geq 1$ ．Therefore，for each $k=1,2, \ldots,\lfloor L\rfloor$ ，we can obtain a doubly covered rectangle of size $p_{0} b$ and $k p_{0} a$ ．

In addition，we have the two ways of folding the ribbon in half along the long axis （leading to a $L \times \frac{1}{2}$ rectangle）or along the short axis（leading to a（ $L / 2$ ）$\times 1$ rectangle）．

We next turn to another idea of folding．Figure 7 a shows how a long ribbon of width 1 can be wrapped by＂winding＂it around a rectangular strip in such a way that the space between successive windings is equal to the width of the ribbon．By bending it backward at the end，as in Figure 7b－c，one obtains a doubly covered strip．Figure 7d
a）

b）


図7 リボンで二重被覆長方形を折る．わかりやすいように，リボンの片方の面に影をつけてある Fig． 7 Folding a ribbon to a doubly－covered rectangle．For better visibility，one side of the ribbon is shaded．
shows the geometric construction：start with a right triangle $A B C$ with the long side $d=B C=\cot \alpha+\tan \alpha$ on a long edge of the ribbon and the right angle $A$ on the opposite edge．When the length $L$ of the ribbon is an even multiple of $d(L=2 n \cdot d)$ ， the folding will close into a doubly covered rectangle．


図8 二重被覆長方形をリボンで折るための別の方法．
Fig． 8 A different way of folding a ribbon to a doubly－covered rectangle

The minimum possible value of $d$ is 2 ．$d$ changes continuously with $\alpha$ ，and any value of $d$ larger than 2 can be obtained．So $n$ ，the number of repetitions，can take all values between 1 and $n_{\max }:=\lfloor L / 4\rfloor$ ．For each $n$ in this range，one can form a right triangle $A B C$ with hypotenuse $d=L /(2 n)$ and legs $\frac{1}{2}\left(\sqrt{d^{2}+2 d} \pm \sqrt{d^{2}-2 d}\right)$ ．One can use the longer leg as the wrapping direction，as in Figure 7，or the shorter leg，as in Figure 8. This leads to doubly covered rectangles of dimensions $\left(n \cdot \frac{1}{2}\left(\sqrt{d^{2}+2 d}+\sqrt{d^{2}-2 d}\right)\right) \times$
$\frac{1}{2}\left(\sqrt{d^{2}+2 d}-\sqrt{d^{2}-2 d}\right)$ and $\frac{1}{2}\left(\sqrt{d^{2}+2 d}+\sqrt{d^{2}-2 d}\right) \times\left(n \cdot \frac{1}{2}\left(\sqrt{d^{2}+2 d}-\sqrt{d^{2}-2 d}\right)\right)$
For $d=2$ ，the two possibilities coincide．So the total number of possibilities is
$\lfloor L / 4\rfloor+\lceil L / 4\rceil-1$ ．This equals $2\lfloor L / 4\rfloor$ except when $L$ is a multiple of 4 ．In this case， we have to subtract 1 to compensate the overcounting for the case $d=2$ ．
But we can see that each doubly covered rectangle by winding can be also obtained by twisting．Hence we obtain $2+\lfloor L\rfloor$ different doubly covered rectangles in total．

## 4．Non－orthogonal polygons that fold to two incongruent boxes

Figure 9 shows a common unfolding of a $4 \times 4 \times 8$ box and a $\sqrt{10} \times 2 \sqrt{10} \times 2 \sqrt{10}$ box．It was obtained by solving an integer programming problem．The integer pro－ gramming model formulates the problem of selecting a subset of 160 unit squares of the axis－aligned square grid underlying Figure 9，subject to the following constraints．
（1）They should form a connected set in the plane．
（2）When folded on the $4 \times 4 \times 8$ box，every square of the surface is covered exactly once．（There are no overlaps．）
（3）When folded on the $\sqrt{10} \times 2 \sqrt{10} \times 2 \sqrt{10}$ box，every part of the surface is covered exactly once．Note that the surface of the $\sqrt{10} \times 2 \sqrt{10} \times 2 \sqrt{10}$ box can be parti－ tioned into 160 unit squares，which are however not aligned with the edges of the box．These squares result from folding the standard grid onto the box surface as shown in Figure 9．Some of these squares bend across an edge of the box．
The algorithm of Section 2 can be viewed as a systematic incremental way of finding all solutions to this problem．
The dimensions of the boxes were chosen as follows：A $1 \times 1 \times 2$ box has surface area 10 ，and a $1 \times 2 \times 2$ box has surface area 16 ．By scaling the first box with the factor 4 and the second box with the factor $\sqrt{10}$ ，we get two boxes with equal surface areas． A square lattice of side length $\sqrt{10}$ can be embedded on the standard integer grid by choosing the vector $\binom{1}{3}$ as a generating＂unit vector＂．
The alignment of the two box unfoldings，with the symmetric layout of two＂central＂ faces sharing two vertices，was fixed and was chosen by hand．
Figure 10 has been made from Figure 9 in an attempt to conceal the obvious folding directions．Further puzzles along these lines（for printing and cutting out）are given on
a web page＊${ }^{\star 1}$ ．


図 9 二つの異なる箱が折れる共通の展開図．ある箱を折るための折り線は，他方の箱を折るための折り線とは直交 しない。交わる角度は $\arctan 3 \approx 72^{\circ}$ である
Fig． 9 A common development of two different boxes．The set of folding lines for one box intersect the other set by neither 90 nor 45 degrees，but at $\arctan 3 \approx 72^{\circ}$ ．

## 5．Concluding remarks

It is an open question if a polygon exists that can fold to three or more orthog－ onal boxes such that all of them have positive volume．We are exploring the pos－ sibility to find such examples by our integer programming model of Section 4 ．If we take the approach in Section 2，the smallest $S$ with $|P(S)| \geq 3$ is given by $P(23)=\{(1,1,11),(1,2,7),(1,3,5)\}$ ．Thus we need to construct polygons of surface area 46 ，which is much bigger than 22 ．

In Section 3．2，we use three different ideas for folding a rectangular ribbon $R$ to a


図 10 二つの異なる箱が折れる共通の展開図．図 9 の境界線を変更して作成。
Fig． 10 A common development of two different boxes．This has been obtained from Figure 9 by modifying the boundary．
doubly－covered rectangle．It would be interesting to classify all ways of folding ribbons into doubly－covered rectangles．In fact，we can generalize the ideas of＂twisting＂and ＂winding＂；see Figure 11．These folding ways correspond to a kind of the billiard ball problem on a rectangular table．Hence，to specify all the folding ways in the figures，we have to find all pairs of relatively prime integers $p$ and $q$ with $p q=\lfloor c L\rfloor$ for $c=1,1 / 4$ ． The number of such pairs seems to be related to the maximal value of prime divisors of numbers in reduced residue system for $\lfloor c L\rfloor^{\star 2}$ ．
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[^3]

図 11 二重被覆長方形を折りための「ねじり折り」と「巻き付け折り」の一般化．
Fig． 11 A generalization of twist／wind folding to a doubly covered rectangle．
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[^0]:    1 Massachusetts Institute of Technology
    $\dagger 2$ 北陸先端科学技術大学院大学
    Japan Advanced Institute of Science and Technology
    $\dagger 3$ Freie Universität Berlin

[^1]:    ＊1 http：／／www．jaist．ac．jp／～uehara／etc／origami／net／all－22．html

[^2]:    $\star 1$ The first program with a naive implementation was too slow．We tuned it with many technical tricks，and now it outputs $L_{22}$ in around 10 hours．

[^3]:    ＊2 http：／／oeis．org／A051265

