

An Efficient Algorithm for Chebyshev Expansion

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1. Introduction

We expand a given function $f(x)$, "well-behaved" over the range $-1 \leq x \leq 1$, by a series of Chebyshev polynomials:

$$f(x) = \sum_{k=0}^{\infty} a_k T_k(x), \tag{1 a}$$

where

$$a_k = \frac{2}{\pi} \int_0^{\pi} (f \cos \theta) \cos k\theta d\theta, \quad T_k(x) \equiv \cos(k \cos^{-1} x). \tag{1 b}$$

A prime indicates that the first term of the series needs to be halved.

Let us consider to evaluate all the coefficients $\{a_k\}$ by numerical quadrature with preassigned accuracy and to reduce the number of operations.

The following two summation rules are known.

$$\frac{2}{n} \sum_{k=1}^n T_r(x_k) T_s(x_k) = \begin{cases} 0, & s=2mn \pm r, \quad r=n \text{ and } s \neq 2mn \pm r \\ (-1)^m, & s=2mn \pm r, \quad r \neq 0, \quad n \\ (-1)^m \times 2, & s=2mn \pm r, \quad r=0 \end{cases} \tag{2}$$

$r=0, 1, \dots, n; \quad m=0, 1, 2, \dots$
 $x_k = \cos((2k-1)/2n)\pi$

$$\frac{2}{n} \sum_{k=0}^n T_r(x_k') T_s(x_k') = \begin{cases} 0, & s \neq 2mn \pm r \\ 1, & s=2mn \pm r, \quad r \neq 0, \quad n \\ 2, & s=2mn \pm r, \quad r=0, \quad n \end{cases} \tag{3}$$

$r=0, 1, \dots, n; \quad m=0, 1, 2, \dots$
 $x_k' = \cos(k/n)\pi$.

A double prime in the above equality means that the first and the last terms are halved.

Following the summation rule (2) and (3), we have classical Chebyshev expansion

$$f(x) = \sum_{r=0}^{n-1} b_r T_r(x), \tag{4 a}$$

$$b_r = \frac{2}{n} \sum_{k=1}^n T_r(x_k) f(x_k), \tag{4 b}$$

and the other Chebyshev expansion

$$f(x) \approx \sum_{r=0}^n c_r T_r(x), \tag{5 a}$$

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$$c_r = \frac{2}{n} \sum_{k=0}^n T_r(x_k') f(x_k') \quad (5 \text{ b})$$

respectively.

It is easily seen that formulas (4b) and (5b) are the approximations of integral (1b) by midpoint rule and by trapezoidal rule, in respect of argument θ . Each coefficient b_r and c_r will be abbreviately called midpoint coefficient and trapezoidal coefficient, respectively.

On the errors of coefficients b_r and c_r , following relations

$$b_r - a_r = -(a_{2n-r} + a_{2n+r}) + (a_{4n-r} + a_{4n+r}) - \dots, \quad (6 \text{ a})$$

$$c_r - a_r = (a_{2n-r} + a_{2n+r}) + (a_{4n-r} + a_{4n+r}) + \dots \quad (6 \text{ b})$$

are known from the two summation rules mentioned above.

Combining these two relations (6a) and (6b), the more accurate coefficients have been found as follows [1]:

$$\frac{c_r + b_r}{2} - a_r = (a_{4n-r} + a_{4n+r}) + \dots,$$

$$\frac{c_r - b_r}{2} - a_{2n-r} = a_{2n+r} + a_{6n-r} + \dots.$$

On the other hand, substituting $2n$ into n in (6b), we find that

$$c_{2n-r}' = a_r + (a_{4n-r} + a_{4n+r}) + \dots,$$

$$c_{2n-r}' = a_{2n-r} + a_{2n+r} + a_{6n-r} + \dots.$$

These trapezoidal coefficients $\{c_r'\}$ determine a truncated Chebyshev series of degree $2n$.

Therefore, we have simple relations between $\{b_r\}$, $\{c_r\}$ and $\{c_r'\}$ such that

$$c_r' = (c_r + b_r)/2 \quad (7 \text{ a})$$

$$c_{2n-r}' = (c_r - b_r)/2, \quad r=0, 1, \dots, n-1 \quad (7 \text{ b})$$

$$c_n' = c_n/2. \quad (7 \text{ c})$$

These relations play an important role in this paper.

2. Evaluation of the coefficients by midpoint rule

Taking $x_k' = \cos(k/2n)\pi$, series (4b) can be rewritten in the form

$$(n/2)b_r = \sum_{k=1}^n T_r(x'_{2k-1}) f(x'_{2k-1}).$$

Neglecting the detailed description [2], we now consider to evaluate the midpoint coefficients $\{b_r\}$. Let n be an even number. We define

$$F_{2k} \equiv f(x'_{2k-1}) + f(-x'_{2k-1}), \quad F_{2k-1} \equiv f(x'_{2k-1}) - f(-x'_{2k-1}),$$

$$F_{4k}^* \equiv F_{2k} + F_{n-2k+2}, \quad F_{4k-2}^* \equiv F_{2k} - F_{n-2k+2},$$

$$F_{4k-1}^* \equiv x'_{2k-1} F_{2k-1} + x'_{n-2k+1} F_{n-2k+1}, \quad F_{4k-3}^* \equiv x'_{2k-1} F_{2k-1} - x'_{n-2k+1} F_{n-2k+1}$$

and

$$\bar{b}_r \equiv (b_{r-1} + b_{r+1})/2, \quad r: \text{even} \quad (8)$$

Then, the following relations are obtained.

$$(n/2)b_r = T_r(x_1') F_{4k}^* + T_r(x_3') F_{4k-2}^* + \dots + T_r(x_{n/2-1}') F_{4k-3}^*, \quad (9)$$

$$r=0, 4, 8, \dots, n-4$$

$$(n/2)b_r = T_r(x_1')F_2^* + T_r(x_3')F_6^* + \cdots + T_r(x'_{n/2-1})F_{n-2}^*,$$

$$r=2, 6, 10, \cdots, n-2 \quad (10)$$

$$(n/2)\bar{b}_r = T_r(x_1')F_3^* + T_r(x_3')F_4^* + \cdots + T_r(x_{n/2-1}')F_{n-1}^*,$$

$$r=0, 4, 8, \cdots, n-4 \quad (11)$$

$$(n/2)\bar{b}_r = T_r(x_1')F_1^* + T_r(x_3')F_5^* + \cdots + T_r(x'_{n/2-1})F_{n-3}^*,$$

$$r=2, 6, 10, \cdots, n-2. \quad (12)$$

These four series with $n/4$ terms can be evaluated by a recurrence formula. If b_r and \bar{b}_r are obtained for $r=0, 2, 4, \cdots, n-2$, the coefficients $b_1, b_3, \cdots, b_{n-1}$ are easily determined by the following formula

$$b_{r+1} = 2\bar{b}_r - b_{r-1}, \quad r=2, 4, \cdots, n-2$$

$$b_1 = \bar{b}_0. \quad (13)$$

As we stated above, all the coefficients $b_r, r=0, 1, \cdots, n-1$ are obtained with $n^2/4$ multiplications approximately, when n is multiple of 4. By slight modification, however, all the midpoint coefficients may be obtained for even number n with the same number of operations.

3. Successive approximation for functions

We now assume that n th degree truncated Chebyshev series has been determined for given function $f(x)$ by the use of trapezoidal rule. Using this n th degree polynomial, we show a method to construct $2n$ th degree polynomial (both polynomials are truncated Chebyshev series with trapezoidal coefficients). Since interpolating points of these polynomials are distributed symmetrically on the range $-1 \leq x \leq 1$ in respect of origin, we concern only on the points on $[0, 1]$.

Let us rewrite the notations as follows. Let $x_k = \cos(k/n)\pi$ be the interpolating point of n th degree polynomial with the trapezoidal coefficients $\{c_k\}$. Similarly, let $x_k' = \cos(k/2n)\pi$ be the interpolating point of $2n$ th degree polynomial with the trapezoidal coefficients $\{c_k'\}$. Then, we have

$$x_{2k}' = x_k, \quad k=0, 1, \cdots, n/2,$$

$$x_1' = \sqrt{(1+x_2')/2}, \quad (14)$$

$$x_{2k-1}' = (x_{2k}' + x_{2k-2}')/(2x_1'), \quad k=2, 3, \cdots, n/2.$$

Given $\{x_k'\}$ and $f(x_{2k-1}')$, $n-1$ th interpolating polynomial is constructed by midpoint rule as stated in paragraph 2. Thus, midpoint coefficients $\{b_r\}$ are obtained. Combining $\{b_r\}$ and $\{c_r\}$ by formula (7), coefficients $\{c_r'\}$, $r=0, 1, \cdots, 2n$ are easily evaluated. Repeating these operations by substituting $2n$ to n , $\{c_r'\}$ to $\{c_r\}$ and $\{x_k'\}$ to $\{x_k\}$, we construct a sequence of interpolating polynomials whose degrees are the power of 2. Setting $n=2$, starting values necessary for successive approximation are given as follows:

$$x_0=1, \quad x_1=1, \quad (15a)$$

and

$$c_0 = \frac{1}{2}f(1) + f(0) + \frac{1}{2}f(-1),$$

$$c_1 = \frac{1}{2}f(1) - \frac{1}{2}f(-1), \quad (15 \text{ b})$$

$$c_2 = \frac{1}{2}f(1) - f(0) + \frac{1}{2}f(-1).$$

If the given function $f(x)$ is sufficiently smooth on $[-1, 1]$, we can determine the stopping rule

$$|c_{n-1}| + |c_n| < \varepsilon \quad (16)$$

for preassigned accuracy ε .

To construct a polynomial sequence mentioned above taking the degree $n=2^1, 2^2, \dots, 2^m$ successively, we require the following operations. Denoting $N=2^m$, we need $N^2/12$ multiplications approximately. Numbers of calculations of square root and of evaluations of the function are $\log_2 N - 1$ and $N+1$, respectively.

4. Conclusion

We can refine the Clenshaw-Curtis method [3], and reduce the number of operations as stated above.

We conjecture that the number of multiplication to calculate the midpoint coefficients for the degree $2^k - 1$ may be more reduced without increasing the complexity of the algorithm, stated in paragraph 2.

References

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