

Minimization of Hazard-Free Switching Networks

TERUO FUKUMURA*, YASUYOSHI INAGAKI*

AND MOTOMU KOMURA**

1. Introduction

With the development of extremely high speed switching networks, malfunctions due to spurious transient outputs, that is, hazards are becoming increasingly critical. Problems of hazards have been discussed by Huffman⁽¹⁾, McCluskey⁽²⁾ and others. They have shown existence condition of hazards and detecting procedure of them.

In this paper we try to design the minimal hazard-free switching networks. At first, the definition of hazards and their existence conditions are given along the lines of McCluskey's work⁽²⁾, and we propose a designing procedure of the minimal hazard-free 2-level AND-OR networks, which is suitable for computer programming. This procedure consists of two steps, i. e., the one is determination of the set of prime implicants satisfying hazard-free conditions and the other is minimal covering. Among them the latter is formalized by the linear programming with 0-1 variables.

Secondly, it is shown that the existence condition of hazards is invariant for the transformation of variables and complementation of variables and functions. Noting this point, for use in synthesis we construct the absolutely minimal hazard-free forms for the representatives of functions of incompletely specified three variables and completely specified four variables. These forms are listed together with invariants in Table 3 and 4. Furthermore we propose the synthesis procedure of the minimal hazard-free networks by the table-look-up scheme using these tables.

2. Design of Minimal Hazard-Free Combinational Networks

2.1. Hazard-Free Condition***

A network whose 1-sets (0-sets) satisfy the following conditions never contains any static or dynamic hazards:

- (1) There are no 1-sets (0-sets) that contain exactly one pair of complementary literals.
- (2) For each pair of adjacent input states that both produce a 1-output (0-output), there is at least one 1-set (0-set) that includes both input states of the pair.

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* Dept. of Electrical Engineering, Faculty of Engineering, Nagoya Univ., Nagoya.

** Formerly a student of graduate school of Nagoya University, now Nippon Telephone Telegraph Corporation, Tokyo.

*** The readers could refer the McCluskey's article⁽²⁾ concerning the definition of hazard and the relating terminology.

2.2. 1-sets Satisfying Condition (2) of Sec. 3.1 and the Procedure Obtaining the Prime Implicants Covering Them

Under the assumption that the input states producing 1-output and don't care input states (unspecified input states) are given using the minimal term representation, we consider the design problem of the minimal hazard-free 2-level AND-OR networks. We can put aside the condition (1) of sec. 2.1 from our consideration because the minimal term which produces 1-output can be a stable 1-set. Therefore we only consider the condition (2) of sec. 2.1. That is, our aim in this section becomes to construct 1-sets which satisfy condition (2) and the prime implicants covering them. In the following we will show a systematic method suitable for digital computers.

Any logical function of n variables $f(x_1, \dots, x_n)$ can be expressed by sum of products expression as follows:

$$f(x_1, x_2, \dots, x_n) = \sum_{\mathbf{e}} f(\mathbf{e}) x_1^{e_1} x_2^{e_2} \dots x_n^{e_n}, \tag{1}$$

where $e_i \in \{0, 1, 2\}$, $\mathbf{e} = (e_1, e_2, \dots, e_n)$

$$x_i^{e_i} = \begin{cases} x_i' (\text{complement of } x) & ; e_i = 0 \\ x_i & ; e_i = 1 \\ 1 & ; e_i = 2 \end{cases}$$

and $f(\mathbf{e})$ is equal to 1 only when the product $x_1^{e_1} x_2^{e_2} \dots x_n^{e_n}$ is contained as a term in the product-sum form.

While a vector \mathbf{e} determines a logical product, it also geometrically determines a vertex, edge, plane, cube etc. of n dimensional hypercube.

[Definition 1] The dimension of cube \mathbf{e} is the number of its coordinates equal to 2. The cube whose dimension is k will be called k -cube.

[Definition 2] For two cubes $\mathbf{a} = (a_1, a_2, \dots, a_n)$ and $\mathbf{b} = (b_1, b_2, \dots, b_n)$, $a_i, b_i \in \{0, 1, 2\}$, we define an operation $\mathbf{a} \circ \mathbf{b}$ as follows:

$$\mathbf{a} \circ \mathbf{b} = (a_1 \circ b_1, \dots, a_n \circ b_n), \tag{2}$$

where $a_i \circ b_i$ is defined as shown in Table 1.

Table 1. \circ operation.

$b_i \backslash a_i$	0	1	2
0	0	y	y
1	y	1	y
2	y	y	2

[Definition 3] The number of y contained in the vector $\mathbf{a} \circ \mathbf{b}$ defined for two cubes \mathbf{a} and \mathbf{b} will be indicated by $\#_y(\mathbf{a} \circ \mathbf{b})$. If $\#_y(\mathbf{a} \circ \mathbf{b}) = 1$, a cube $\mathbf{a} * \mathbf{b}$ will be defined by changing y of $\mathbf{a} \circ \mathbf{b}$ to 2. If $\#_y(\mathbf{a} \circ \mathbf{b}) \geq 2$, cube $\mathbf{a} * \mathbf{b}$ does not exist.

[Definition 4] Cube \mathbf{a} is said to include \mathbf{b} when the condition of either $a_i = b_i$ or $a_i = 2$ whenever $a_i \neq b_i$ is satisfied.

We note from the above definitions that cube $\mathbf{a} * \mathbf{b}$ is the one that includes both of \mathbf{a}

and \mathbf{b} and whose dimension is greater than them by 1.

Therefore we can get all the desired 1-sets and prime implicants according to the following procedures A and B.

Procedure A Procedure of obtaining 1-sets which satisfy condition (2) of sec. 2.1.

- (1) Represent the given input states which produce 1-output by the vectors \mathbf{e} used in Eq. (1). The vectors obtained are 0-cubes.
- (2) Find all 1-cubes $\mathbf{a}*\mathbf{b}$ for all pairs of cubes \mathbf{a}, \mathbf{b} . They are the 1-sets which cover the adjacent input states.
- (3) List up all 1-cubes obtained in (2) and 0-cubes which are not included by any of them.

Procedure B Procedure of obtaining prime implicants

- (1) Represent the given input states which produce 1-output and don't care input states by vector \mathbf{e} used in Eq. (1).
- (2) Begin with $k=0$ and find $(k+1)$ -cubes by doing $*$ operation for all pairs of k -cubes. Proceed this procedure till the number of $(k+1)$ -cubes newly obtained becomes zero.
- (3) List up all k -cubes which are not included by any of $(k+1)$ -cubes.

2.3. Procedure for Designing Minimal Hazard-Free Networks

To construct a minimal hazard-free network can be reduced to the minimal covering problem finding the minimal set of the prime implicants in the set of the prime implicants obtained by the procedure B, which covers 0-cubes and 1-cubes obtained by the procedure A. Furthermore this minimal covering problem can be formalized by 0-1 variables linear programming.

2.3.1. Formulation by the Linear Programming

Let the 0-cubes and 1-cubes obtained by the procedure A be c_1, c_2, \dots, c_m and let the prime implicants cubes obtained by the procedure B be r_1, r_2, \dots, r_n . Next define the constants $a_{ij} (i=1, \dots, m; j=1, \dots, n)$ as

$$a_{ij} = \begin{cases} 1, & \text{if } r_j \text{ includes } c_i, \\ 0, & \text{otherwise,} \end{cases}$$

and also the 2-valued variables $z_j (j=1, \dots, n)$ as

$$z_j = \begin{cases} 1, & \text{if } r_j \text{ is contained in the desired set of the prime implicants,} \\ 0, & \text{otherwise,} \end{cases}$$

then our problem can be formulated as follows. That is,

Procedure C Under the constraints

$$\sum_{j=1}^n a_{ij} z_j \geq 1 \quad (i=1, \dots, m) \quad (3)$$

find $z_j (=0 \text{ or } 1, j=1, \dots, n)$ minimizing the objective functions,

$$(1) \quad y = 1 + \sum_{j=1}^n z_j \quad (\text{total number of AND and OR gates}) \quad (4)$$

or

$$(2) \quad y = \sum_{j=1}^n w_j z_j + \sum_{j=1}^n z_j \quad (\text{total number of diodes}), \quad (5)$$

where $w_j = N - d_j$, d_j is the dimension of cube r_j and N is the number of input variables.

2.3.2. Reduction of Number of the Constraints and Variables

Many algorithms of the integer linear programming have been proposed and, for examples, Gomory's⁽³⁾ and Balas'⁽⁴⁾ methods can be employed to execute the procedure C. However, in any execution it is of course desirable to decrease the number of constraints and variables. For this purpose we utilize the following three properties.

(Property 1) If $a_{ij_1} = 1$ and $a_{ij} = 0 (j \neq j_1)$ then $z_{j_1} = 1$. Therefore, the variable z_{j_1} and the constraints such that $a_{kj_1} = 1$ can be put aside from our consideration.

(Property 2) If for the i_1 -th and i_2 -th constraints, $a_{i_2j} = 1$ whenever $a_{i_1j} = 1$, the i_2 -th constraint can be omitted.

(Property 3) If for the coefficients of z_{j_1} and z_{j_2} , $a_{i_2j} = 1$ whenever $a_{i_1j} = 1$, z_{j_1} is equal to 0.

These properties (1), (2) and (3) respectively correspond to (1) essential term, (2) column dominance, and (3) row dominance found in the prime implicant table which is obtained from the $m \times n$ matrix (a_{ij}) by replacing 1-elements by \vee symbols and removing 0-elements.

3. The Same Class Transformation and Number of Classes of Functions

3.1. The Same Class Transformation and Hazard

The transformation consisting of (1) complementation of variable, (2) permutation of variables and (3) complementation of function will be called the same class transformation of Boolean function. The relation among the functions which are convertible to each other by finite use of operations of the same class transformation is indicated by R .

This relation R is an equivalence relation on the family of Boolean functions and thus the functions are resolved into equivalence classes. The functions which are involved in a same class will be called the same class function.

An incompletely specified n variables Boolean function is a many to one mapping from $\{0, 1\}^n$ (n order direct product of the space $\{0, 1\}$) to the space $\{0, 1, d\}$. In other words, the function assigns one of 0, 1 and d to each of 2^n vertices in n dimensional hypercube.

Considering the fact that permutation and negation of variables, which involved in the same class transformation, are irrelevant to the Hamming distances between vertices, we can say that if two figures whose vertices are assigned 0, 1 and d (a part of n dimensional hypercube) are congruent, they correspond to the same class functions.

The same class transformation is nothing but a combination of relabeling transformations of variables and transformations between 0 and 1. Permutation of variables changes the labels of elements in 1-sets (0-sets), while complementation of variables yields complementation of the corresponding elements in 1-sets (0-sets) and complementation of function makes the elements of 1-set (0-set) be complemented and become the element of 0-set (1-set). Therefore, neither permutation nor complementation evidently produce a new hazard. Furthermore, from the fact that 1-sets and 0-sets have mutually dual relations it is also clear that any new hazard is not yielded even by complementation of function.

Thus we get the following assertion concerning the effect of the same class transformation on hazard.

[Assertion 1] There is no possibility of any hazard in any network which is derived from a hazard-free network by using the same class transformation.

3.2. Enumeration of Incompletely Specified Boolean Functions

It is natural to ask for the number of classes of incompletely specified Boolean functions in connection with their design algorithm. Many results and techniques⁽⁵⁾ concerning the problem for completely specified functions have been already obtained by using the group theory. The same techniques can be extended to the case of incompletely specified functions, yielding more general results as is expected. The latter results include the former as the results of the special case.

Numbers of classes of incompletely and completely specified functions are tabulated in Table 2 for several values of n .

Table 2. Number of classes.

n	completely specified function*			incompletely specified function		
	total	complement, permutation of variable	containing complement of function	total	complement, permutation of variable	containing complement of function
2	16	6	4	81	21	13
3	256	22	14	6,561	267	155
4	65,536	402	238	43,046,721	132,102	about 60,000
5	4,294,967,296	1,228,158	about 600,000	—	—	—

* the results for the completely specified function are taken from [5].

3.3. Invariant of Incompletely Specified Function and Representative Function

The invariants which characterize the class of functions will be computed in this section. The weight of a function f will be denoted by a pair of $w_1(f)$ and $w_d(f)$, the former represents the number of 1-vertices of a function and the latter stands for the sum of numbers of 1-vertices and d -vertices of the function.

Given a function f of n variables, its invariants are computed as follows:

[Algorithm of computing the invariants]

(1) Compute the sequence of the weights

$$\{w_d(f), w_d(f \oplus x_1), \dots, w_d(f \oplus x_n), w_d(f \oplus x_1 \oplus x_2), \dots, w_d(f \oplus x_{n-1} \oplus x_n), \dots, w_d(f \oplus x_1 \oplus \dots \oplus x_n), w_1(f), w_1(f \oplus x_1), \dots, w_1(f \oplus x_1 \oplus \dots \oplus x_n)\},$$

where \oplus denotes the ring sum (exclusive or).

(2) If $w_1(f) + w_d(f) > 2^{n-1}$, convert 0-vertices into 1-vertices and vice versa, where d -vertices are not changed, and compute again the sequence of weights. This procedure corresponds to the complementation of the function.

(3) If $w_d(f \oplus x_i) > 2^{n-1}$, complement all those weights which involve x_i , i. e., w goes to $2^n - w$. When $w_d(f \oplus x_i) = 2^{n-1}$, if complementation of all the weights involving x_i makes the weight sequence smaller than ever in lexicographic order, do this operation. This procedure corresponds to complementation of variables.

(4) Permute $w_d(f \oplus x_i)$ and $w_d(f \oplus x_j)$ so that they are in ascending order. The permutation of the relevant variables must be consistently applied to all the other weights that involve these variables.

If some of the first-order weights $w_d(f \oplus x_i)$ are the same, try all permutations of these weights to obtain the weight sequence which is smallest in lexicographic order.

This procedure corresponds to the permutation of variables.

(5) After the smallest sequence is obtained concerning w_d , apply the procedure (3) and (4) to the sequence of w_1 as far as w_d is unchanged.

The final sequence obtained in this fashion is the smallest sequence of weights and will be called the invariant of a function f . A function f whose sequence of weights is the smallest sequence is said to be of the normal form. We will choose the functions of normal form as the representative functions.

4. *Representative Functions and Synthesis of the Minimal Hazard-Free Networks Using Invariants*

4.1. *Synthesis Procedure*

Given a table in which the minimal hazard-free networks for the representative functions of n variables and their invariants are tabulated, the minimal hazard-free network of any function f of n variables can be obtained as follows:

- (1) Compute the sequence of weights of a given function f .
- (2) Compute the invariants of f using the transformation described in sec. 3.3, and register them sequentially.
- (3) Find the function with the same invariants in the table.
- (4) Realize f by applying the inverse of the transformations registered in the step (2) to the network given in the table.

4.2. *The Minimal Hazard-Free Networks for the Representative Functions*

Choosing the functions of normal forms as the representatives we calculated their invariants and constructed the minimal hazard-free network by use of the techniques described in secs. 2, 3 and 4. This computation was done by NEAC 2203 computer. The results are shown in Tables 3 and 4. That is, all the representatives of incompletely specified 3 variables functions and of completely specified 4 variables functions are tabulated in Tables 3 and 4, respectively.

In the case of completely specified 4 variables functions, the results were obtained without using the linear programming. It took about 2 hours to obtain the invariants and the minimal networks of 238 functions of 4 variables because of the low performance rate of the computer used.

Table 3. Invariants for incompletely specified 3 variables functions and their minimal hazard-free forms.

No.	Invariant	1-vertex	d-vertex	Function	No.	Invariant	1-vertex	d-vertex	Function	
1	04444444	04444444	-	0	81	42244264	22246444	67	34	XY
2	13335553	04444444	-	7	82	42244264	22444462	47	36	YZ+XZ, XY+XZ
3	22246444	04444444	-	67	83	42442246	24244246	36	45	XYZ+XZ
4	24442264	04444444	-	56	84	42442246	24442264	34	56	XYZ+XZ, XY+YZ+XY
5	24442264	04444444	-	34	85	42442246	22464244	46	35	XZ
6	51335535	04444444	-	567	86	42442246	22444266	56	34	XY+XZ
7	33333373	04444444	-	347	87	51333355	22444426	56	347	X
8	33333337	04444444	-	356	88	51333355	24244246	36	457	X+YZ
9	40444444	04444444	-	4567	89	51333355	24224464	37	456	YZ
10	44440444	04444444	-	2345	90	51333355	22246444	67	345	X
11	44444440	04444444	-	1247	91	51333355	22444462	47	356	X
12	42244446	04444444	-	3567	92	51333355	22464244	46	357	X
13	42244264	04444444	-	3467	93	51333355	24442264	34	567	X+YZ
14	42442246	04444444	-	3456	94	53331553	22424644	57	234	XZ
15	53335557	04444444	-	03567	95	53331553	22444462	47	236	XZ+XY, XY+YZ
16	53331553	04444444	-	23457	96	53331553	24422446	35	247	XY+XY, XY+XZ, YZ+XZ, XY+YZ
17	51333355	04444444	-	34567	97	53331553	24442264	34	257	XY+XY, XZ+XY
18	64442224	04444444	-	123456	98	53331553	22642444	45	237	XY
19	62442442	04444444	-	124567	99	53331553	24462442	24	357	XY+XY
20	62242444	04444444	-	234567	100	53335557	24466446	06	357	XY+YZ+XY
21	73333335	04444444	-	1234567	101	53335557	22444426	56	037	XY+XZ
22	84444444	04444444	-	01234567	102	53335557	24446664	07	356	XY+YZ+XY, XY+XZ+XY, XY+XZ+XZ
23	13335553	13335553	7	-	103	53335557	22246444	67	035	XY
24	22246444	13335553	7	6	104	64442224	22444426	56	1234	XZ+XY, YZ+XY
25	22444426	13353335	6	5	105	64442224	22464244	46	135	XZ
26	24442264	13533335	4	57	106	64442224	24442264	34	1256	XY+XY, XY+XZ, XZ+XZ, XZ+XY
27	31335535	13335553	6	57	107	62444422	26444422	12	4567	YZ+XY
28	31335535	13335553	7	56	108	62444422	24244642	27	1456	YZ+X
29	33333373	13335553	7	34	109	62444422	24264424	26	1457	YZ
30	33333373	13533335	4	37	110	62444422	24442624	25	1467	YZ+X
31	33333373	15333355	3	47	111	62444422	22246444	67	1245	X
32	33333337	13353335	6	35	112	62444422	22444426	56	1247	X
33	40444444	13335553	7	456	113	62444422	22444462	47	1256	X
34	44440444	13533335	5	234	114	62242444	22642444	45	2367	X
35	44444440	13335553	7	124	115	62242444	24422442	35	2467	X+Y
36	42224446	13353335	6	357	116	62242444	24442264	34	2567	X+Y
37	42224446	13335553	7	356	117	62242444	22424644	57	2346	X
38	42244264	13335553	7	346	118	62242444	22444426	56	2347	X
39	42244264	13533335	4	367	119	62242444	22246444	67	2345	X
40	42442244	13353335	6	345	120	31335535	31335535	567	-	XY+XZ
41	42442244	15333355	3	456	121	33333373	33333373	347	-	XY+XZ
42	42442244	13533335	4	356	122	33333337	33333337	356	-	XY+XZ+YZ
43	53335557	13335553	7	0356	123	40444444	31335535	567	4	X
44	53335557	13353335	6	0357	124	44440444	33531533	245	3	XY+XY
45	53335557	15353335	0	3567	125	44444440	33333351	247	1	XY+XZ+XY+XZ+XY+YZ
46	53331553	13335553	7	2345	126	42224446	31335535	567	3	XY+XZ
47	53331553	13533335	5	2347	127	42224446	33333337	356	7	YZ+XY+XZ
48	53331553	13533335	4	2357	128	42244264	33333373	347	6	YZ+XZ
49	51333355	13335553	7	3466	129	42244264	31353335	467	3	XY+XZ
50	51333355	13533335	6	3457	130	42442246	33353155	346	5	XYZ+XZ
51	51333355	13533335	4	3567	131	42442246	31553335	456	3	XY+XZ
52	51333355	15333355	3	4567	132	42442246	33333337	356	4	XYZ+XZ+XY
53	64442224	13353335	6	12345	133	51333355	33333373	347	56	X+YZ
54	62444422	13335553	7	12456	134	51333355	31353335	367	45	X+YZ
55	62444422	13533335	6	12457	135	51333355	31533335	456	37	X
56	62444422	15353335	2	14567	136	51333355	33353155	346	57	X+YZ
57	62242444	13335553	7	23456	137	51333355	33333337	356	47	X+YZ
58	62244444	13533335	5	23467	138	51333355	33135355	367	45	X+YZ
59	73333333	13335553	7	123456	139	51333355	31335535	567	34	X
60	73333333	13533335	6	123457	140	53331553	33531533	245	37	XY+XY
61	73333333	13533355	4	123567	141	53331553	33531355	345	27	XY+XY, XY+YZ
62	22246444	22246444	67	-	142	53331553	33353551	247	35	XY+XY+YZ, XY+XY+XZ
63	22444426	22444426	56	-	143	53331553	33333373	347	25	YZ+XY
64	24442264	24442264	34	-	144	53331553	31533553	457	23	XY+XZ
65	31335535	22246444	67	5	145	53331553	33313555	357	24	YZ+XZ
66	31335535	22444426	56	7	146	53335557	31335535	567	03	XY+XZ
67	33333373	24442264	34	7	147	53335557	33357555	067	35	XY+XY+XZ
68	33333373	24224446	37	4	148	53335557	33333337	356	07	YZ+XY+XZ
69	33333373	24444462	47	3	149	53335557	33535537	056	37	XY+XZ+XY+XZ
70	33333337	22444426	56	3	150	40444444	40444444	4567	-	X
71	40444444	22246444	67	45	151	44440444	44440444	2345	-	XY+XY
72	40444444	22444426	56	47	152	44444440	44444440	1247	-	XY+XY+YZ, XY+XY+XZ
73	44440444	26242444	23	45	153	42224446	42224446	3567	-	XY+YZ+XZ
74	44440444	24442264	34	25	154	42244264	42244264	3467	-	XY+YZ+XZ
75	44440444	24422446	35	24	155	42442246	42442246	3456	-	XY+YZ+XZ
76	44444440	24444422	47	12						
77	42224446	22246444	67	35						
78	42224446	22444426	56	37						
79	42244264	22444426	37	46						
80	42244264	24442264	34	67						

(Note) 1. 1-vertexes and 0-vertexes are represented as decimal numbers corresponding to binary numbers which represent vertexes.
 2. W* is the complement of W, XYZ is the logical product of X, Y and Z, X+Y is the logical sum of X and Y.

Table 4. Invariants for 4 variables functions and their minimal hazard-free forms.

No.	Invariant	1-vertex	Decimal representation of function	Function
1	0	0	0	0
2	1	1	40	XYZ
3	2	1 1	41 39	WXZ+XYZ
4	3	1 1 1	37 36	WYZ+XYZ
5	4	1 1 1 1	35 34	WXYZ+XYZ
6	5	1 1 1 1 1	33 32	WXYZ+WXYZ
7	6	1 1 1 1 1 1	31 30	WXYZ+WXYZ+XYZ
8	7	1 1 1 1 1 1 1	29 28	WXYZ+WXYZ+XYZ
9	8	1 1 1 1 1 1 1 1	27 26	WXYZ+WXYZ+XYZ
10	9	1 1 1 1 1 1 1 1 1	25 24	WXYZ+WXYZ+XYZ
11	10	1 1 1 1 1 1 1 1 1 1	23 22	WXYZ+WXYZ+XYZ
12	11	1 1 1 1 1 1 1 1 1 1 1	21 20	WXYZ+WXYZ+XYZ
13	12	1 1 1 1 1 1 1 1 1 1 1 1	19 18	WXYZ+WXYZ+XYZ
14	13	1 1 1 1 1 1 1 1 1 1 1 1 1	17 16	WXYZ+WXYZ+XYZ
15	14	1 1 1 1 1 1 1 1 1 1 1 1 1 1	15 14	WXYZ+WXYZ+XYZ
16	15	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	13 12	WXYZ+WXYZ+XYZ
17	16	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	11 10	WXYZ+WXYZ+XYZ
18	17	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	9 8	WXYZ+WXYZ+XYZ
19	18	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	7 6	WXYZ+WXYZ+XYZ
20	19	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	5 4	WXYZ+WXYZ+XYZ
21	20	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	3 2	WXYZ+WXYZ+XYZ
22	21	1 1	1 0	WXYZ+WXYZ+XYZ
23	22	1 1	0 0	WXYZ+WXYZ+XYZ
24	23	1 1	0 0	WXYZ+WXYZ+XYZ
25	24	1 1	0 0	WXYZ+WXYZ+XYZ
26	25	1 1	0 0	WXYZ+WXYZ+XYZ
27	26	1 1	0 0	WXYZ+WXYZ+XYZ
28	27	1 1	0 0	WXYZ+WXYZ+XYZ
29	28	1 1	0 0	WXYZ+WXYZ+XYZ
30	29	1 1	0 0	WXYZ+WXYZ+XYZ
31	30	1 1	0 0	WXYZ+WXYZ+XYZ
32	31	1 1	0 0	WXYZ+WXYZ+XYZ
33	32	1 1	0 0	WXYZ+WXYZ+XYZ
34	33	1 1	0 0	WXYZ+WXYZ+XYZ
35	34	1 1	0 0	WXYZ+WXYZ+XYZ
36	35	1 1	0 0	WXYZ+WXYZ+XYZ
37	36	1 1	0 0	WXYZ+WXYZ+XYZ
38	37	1 1	0 0	WXYZ+WXYZ+XYZ
39	38	1 1	0 0	WXYZ+WXYZ+XYZ
40	39	1 1	0 0	WXYZ+WXYZ+XYZ
41	40	1 1	0 0	WXYZ+WXYZ+XYZ
42	41	1 1	0 0	WXYZ+WXYZ+XYZ
43	42	1 1	0 0	WXYZ+WXYZ+XYZ
44	43	1 1	0 0	WXYZ+WXYZ+XYZ
45	44	1 1	0 0	WXYZ+WXYZ+XYZ
46	45	1 1	0 0	WXYZ+WXYZ+XYZ
47	46	1 1	0 0	WXYZ+WXYZ+XYZ
48	47	1 1	0 0	WXYZ+WXYZ+XYZ
49	48	1 1	0 0	WXYZ+WXYZ+XYZ
50	49	1 1	0 0	WXYZ+WXYZ+XYZ
51	50	1 1	0 0	WXYZ+WXYZ+XYZ
52	51	1 1	0 0	WXYZ+WXYZ+XYZ
53	52	1 1	0 0	WXYZ+WXYZ+XYZ
54	53	1 1	0 0	WXYZ+WXYZ+XYZ
55	54	1 1	0 0	WXYZ+WXYZ+XYZ
56	55	1 1	0 0	WXYZ+WXYZ+XYZ
57	56	1 1	0 0	WXYZ+WXYZ+XYZ
58	57	1 1	0 0	WXYZ+WXYZ+XYZ
59	58	1 1	0 0	WXYZ+WXYZ+XYZ
60	59	1 1	0 0	WXYZ+WXYZ+XYZ

5. Conclusion

We have led a design procedure of the minimal hazard-free networks of 2-level AND-OR gates using the linear programming techniques. The networks thus obtained are minimal in the sense that total number of AND, OR gates or diodes to be needed are minimum. Next we have shown that the networks obtained by applying the same class transformation to a hazard-free networks do not have any hazard. The existing techniques for calculating the number of classes of functions and determining the representative functions have been extended to the case of incompletely specified functions.

By use of the techniques proposed by us, the minimal hazard-free forms of the representatives of incompletely specified functions of 3 variables and that of completely specified functions of 4 variables together with all their invariants have been obtained and are tabulated in Tables 3 and 4. Finally we have shown the way to use these tables.

We have left a problem to improve the method applicable to the functions of five or more variables. We have also left the problem to find the easy and effective algorithm to solve the 0-1 variables linear programming. But we will have opportunities to attack these problems some other time.

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